Resilient Distributed Datasets, Spark

DS 5110: Big Data Systems (Spring 2023)

Lecture 4

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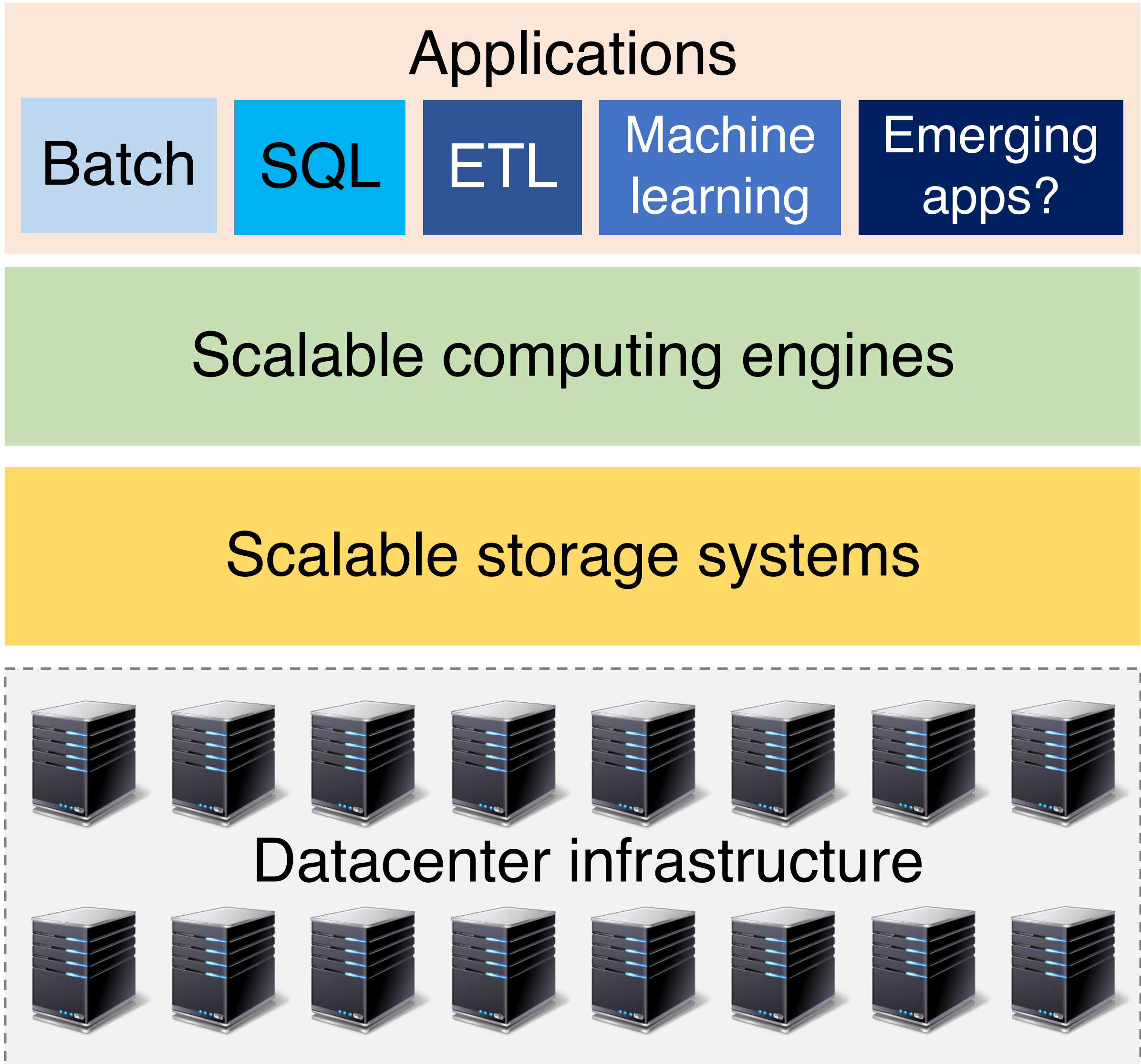
University of Virginia

Some material taken/derived from:

- Matei Zarahia's NSDI'12 talk slides.


Applications
Batch
SQL
ETL
Machine learning
Emerging apps?
Scalable computing engines
Scalable storage systems
Datacenter infrastructure

What's good with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault tolerance as a concern

Problems with MapReduce

- Scaled analytics to thousands of machines
- Eliminated fault tolerance as a concern

- Not very expressive
  - Iterative algorithms
    (PageRank, Logistic Regression, …)
  - Interactive and ad-hoc queries
    (Interactive Log Debugging)

- Lots of specialized frameworks
  - Pregel, GraphLab, PowerGraph, DryadLINQ, HaLoop…

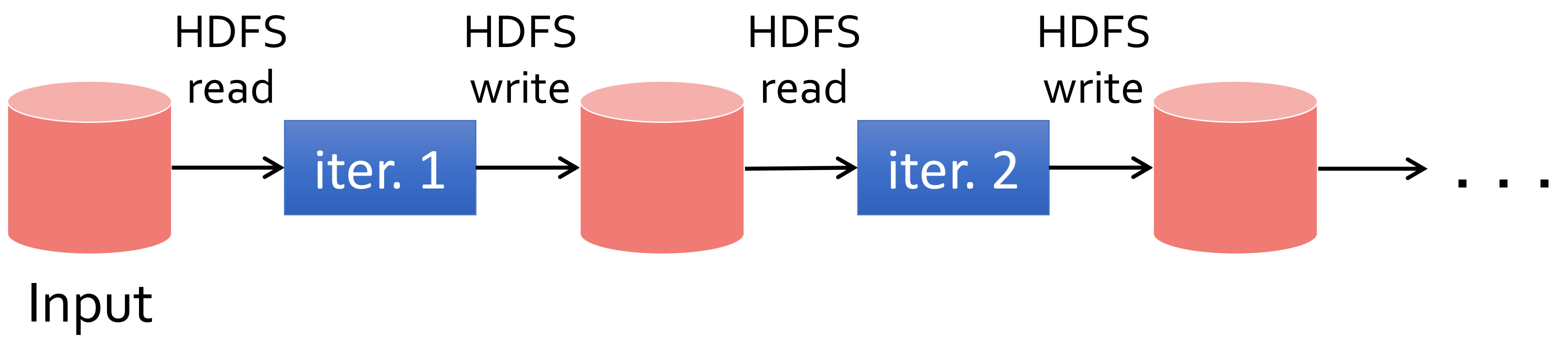
Sharing data between stages/iterations

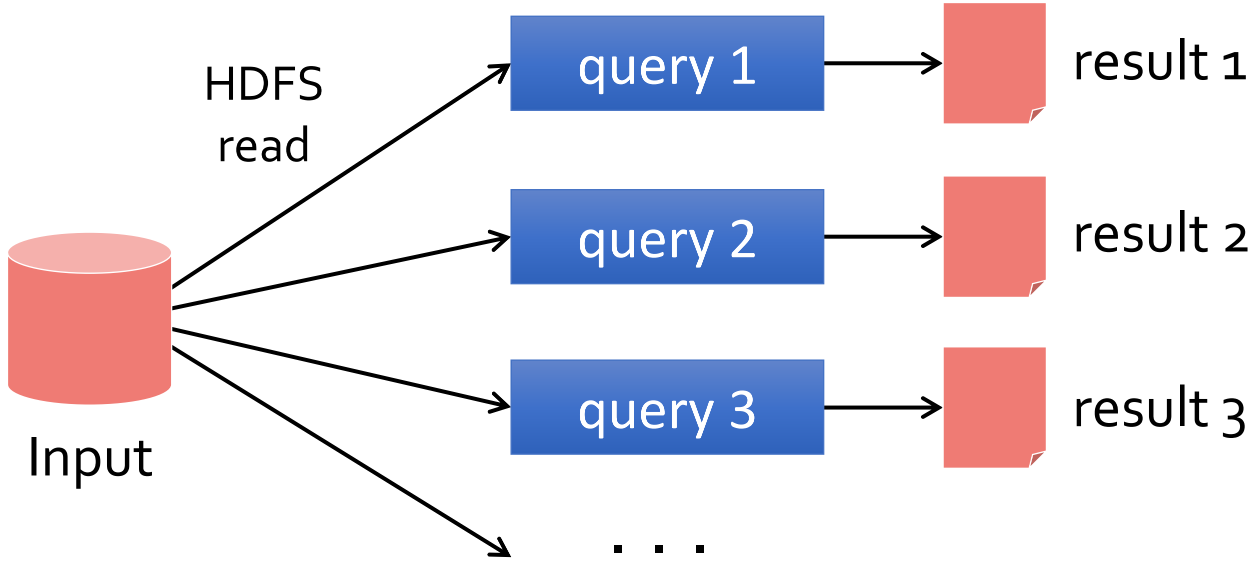
- Only way to share data between iterations / phases is through shared storage
  - Slow!
- Allow operations to feed data to one another
  - Ideally, through memory instead of disk-based storage

Sharing data between stages/iterations

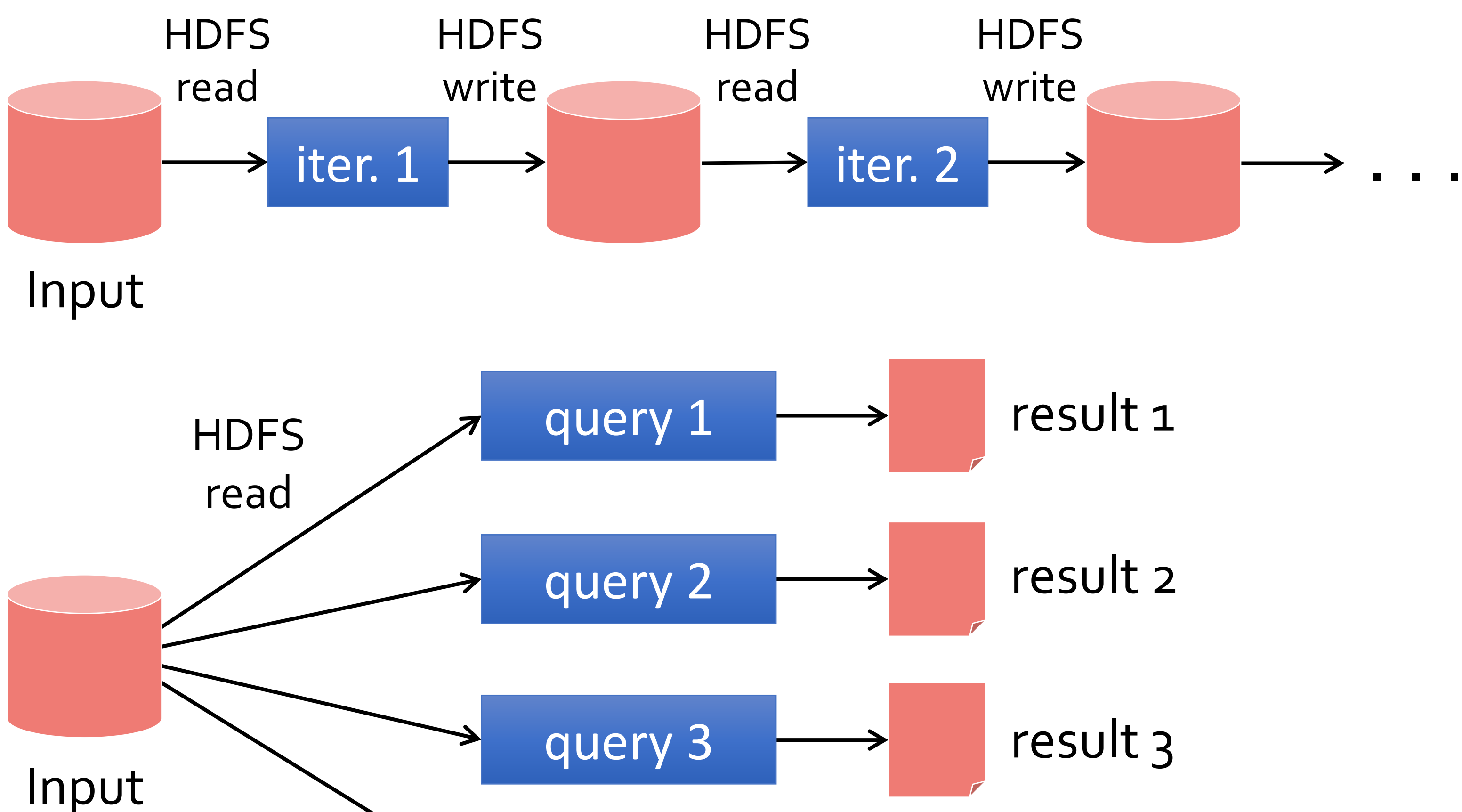
- Only way to share data between iterations / phases is through shared storage
  - Slow!
- Allow operations to feed data to one another
  - Ideally, through memory instead of disk-based storage
- Need the "chain" of operations to be exposed to make this work
- **Problem to solve:** Would this break the MR fault-tolerance scheme?
  - Retry and Map or Reduce task since idempotent

Examples


HDFS
read
HDFS
write
HDFS
read
HDFS
write
iter. 1
iter. 2
. . .
Input

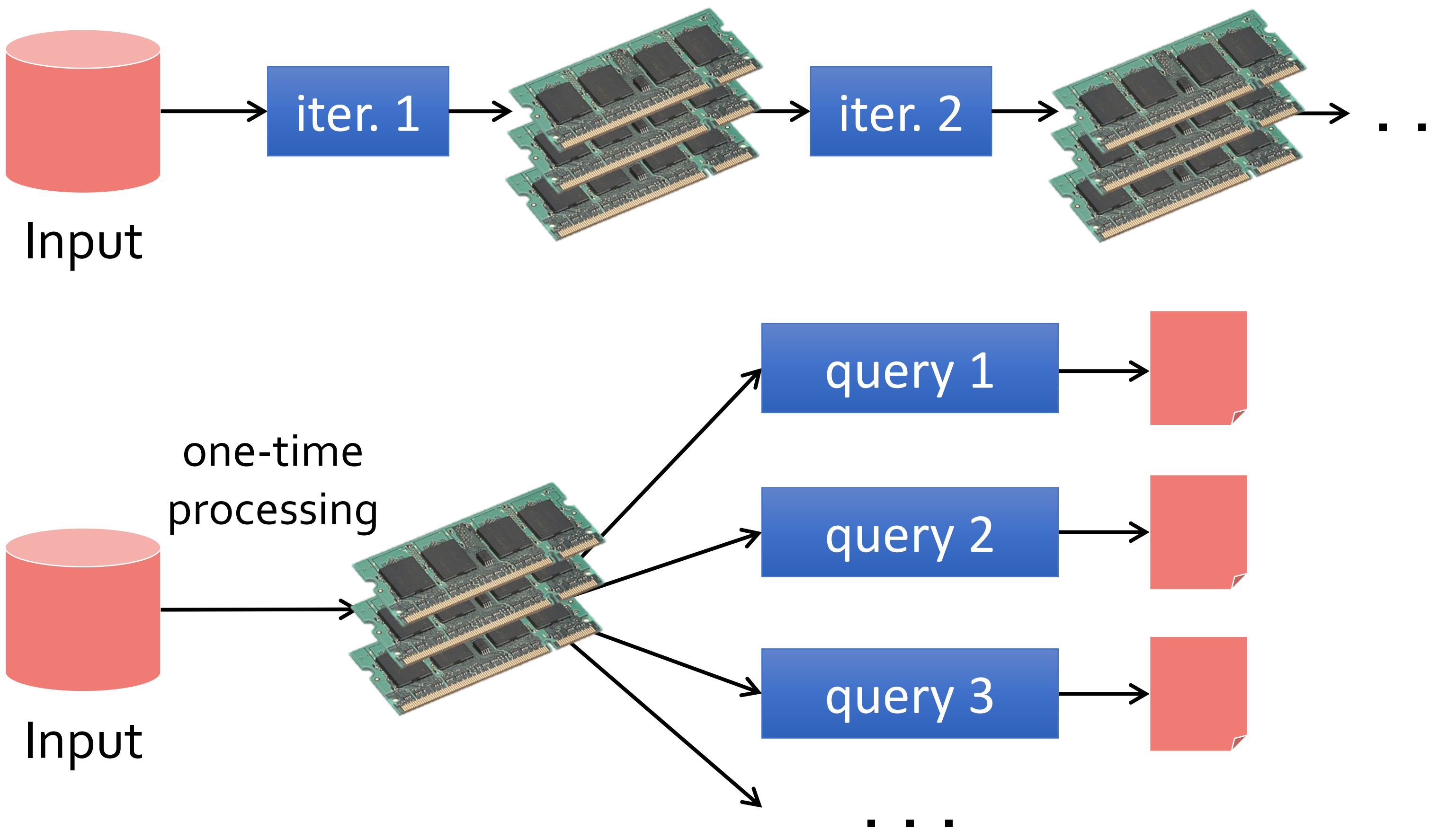

HDFS
read
query 1
query 2
query 3
result 1
result 2
result 3
Input
. . .

Examples



Slow due to replication and disk I/O,
but necessary for fault tolerance

Goal: In-memory data sharing



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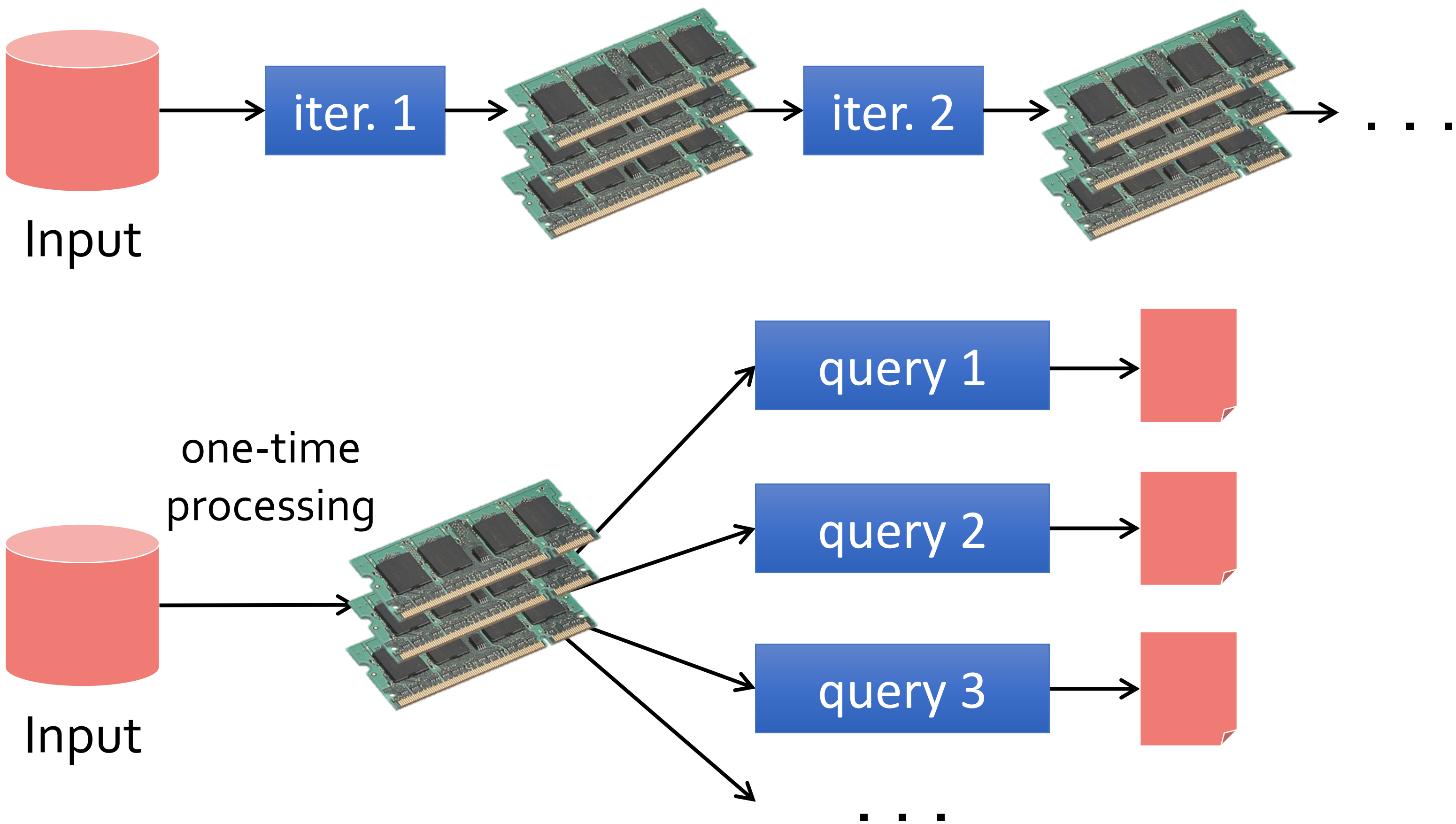

10-100× faster than network/disk, **but how to get FT?**

Challenges

- How to design a distributed memory abstraction that is both fault-tolerant and efficient?

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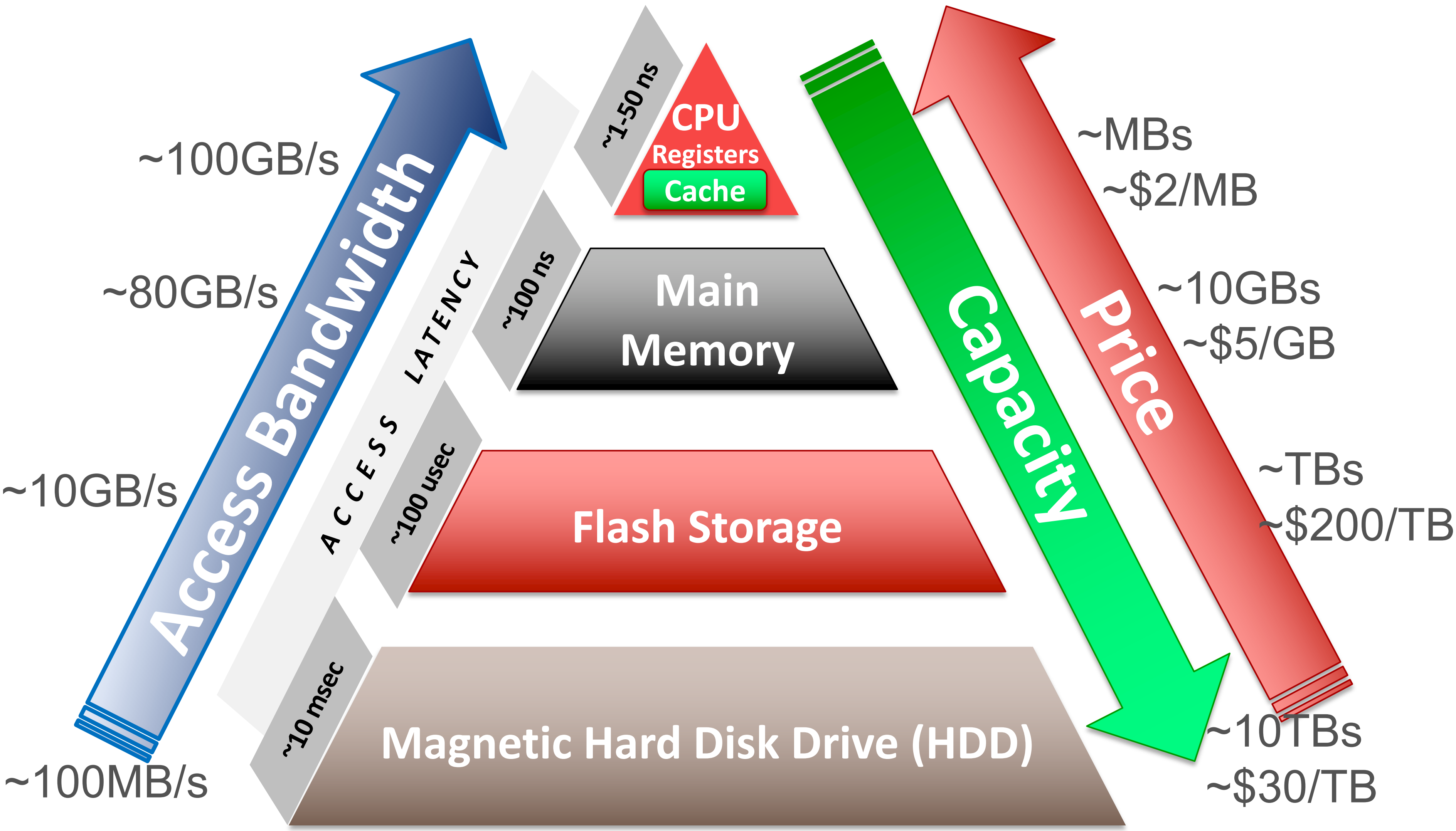
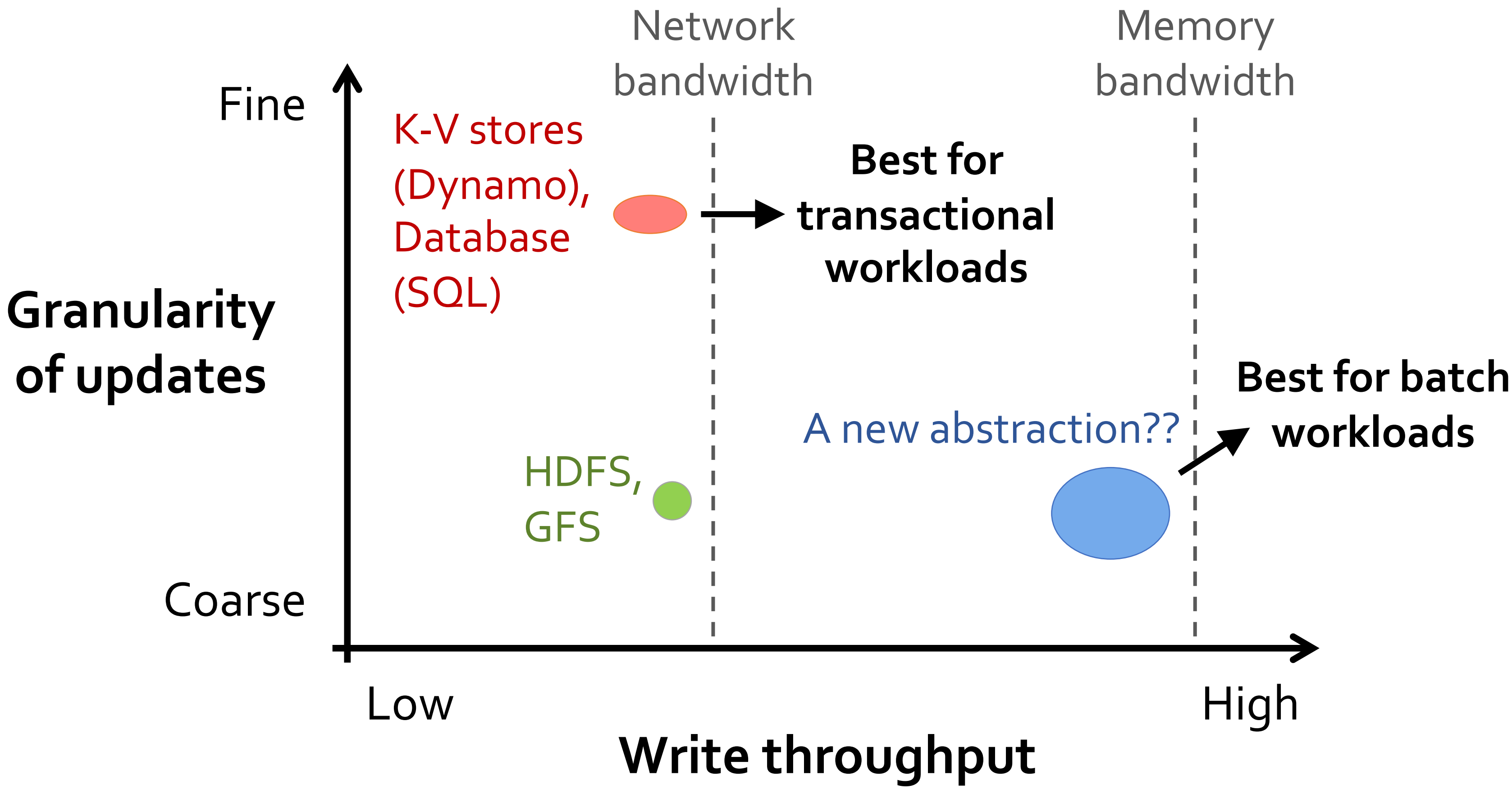
- Existing storage systems allow fine-grained mutation to state
  - In-memory key-value stores
  - Requires replicating data or logs across nodes for fault tolerance
    - Costly for data-intensive apps
    - 10-100x slower than memory write
  - They also require costly on-the-fly replication for mutations

Memory-storage hierarchy



*UCSD DSC 102: Systems for scalable analysis. Arun Kumar

Tradeoff space



Challenges

- How to design a distributed memory abstraction that is both fault-tolerant and efficient?

- Existing storage systems allow fine-grained mutation to state

Insight: leverage similar coarse-grained approach that **transforms whole dataset per operation**, like MapReduce (batch processing)

    - 10-100x slower than memory write
  - They also require costly on-the-fly replication for mutations

Solution: Resilient Distributed Datasets (RDDs)

- Restricted form of distributed shared memory
  - Immutable, partitioned collections of records
  - Can only be built through *coarse-grained,* deterministic ***transformations*** (map, filter, join, …)
- Efficient fault recovery using *lineage*
  - Log one operation to apply to many elements
  - Recompute lost partitions on failure
  - No cost if nothing fails

Spark programming interface

Scala API, now have multi-language bindings such as Python, Java, etc.

Managing RDDs

- Transformations on RDDs ($RDD_1 \rightarrow RDD_2$)
- Actions on RDDs (RDD → output)
- Control over RDD partitioning (how items are split over nodes)
- Control over RDD persistence (in memory, on disk, or recompute on loss)

Transformations

Transformations (define a new RDD)	map filter sample groupByKey reduceByKey sortByKey	flatMap union join cogroup cross mapValues

RDDs in terms of Scala types → Scala semantics at workers

Transformations are lazy operations; cause no cluster action

Actions

Actions (return a result to driver program)	collect reduce count save lookupKey

Consumes an RDD to produce output

either to storage (save), or

to interpreter/Scala (count, collect, reduce)

Causes RDD lineage chain to get executed on the cluster to produce the output

(for any missing pieces of the computation)

Interactive debugging

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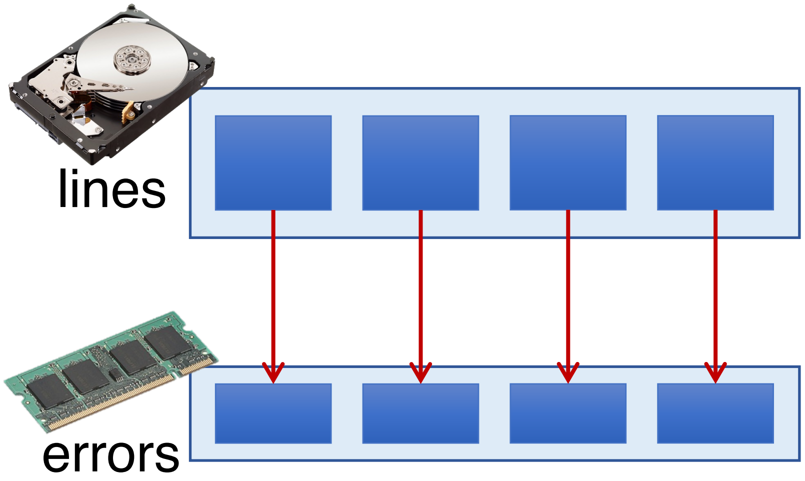

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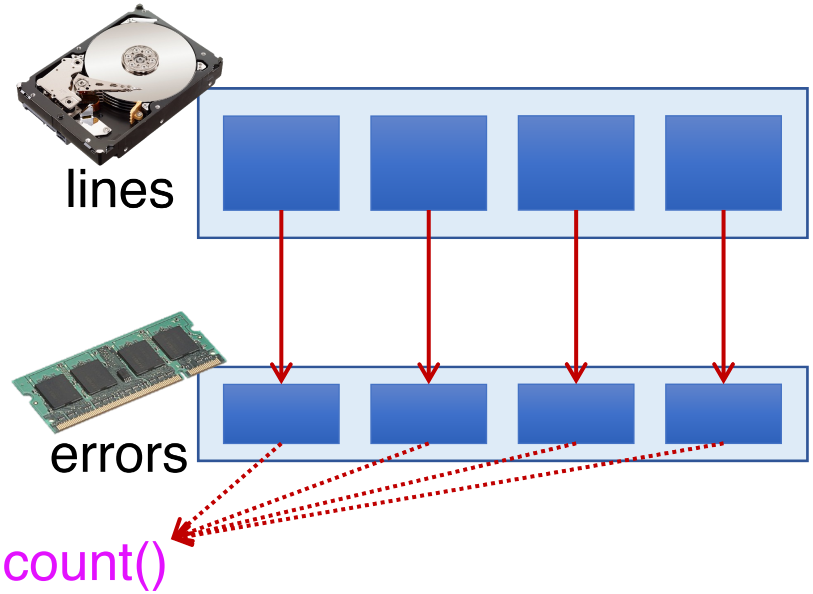

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errors.filter(
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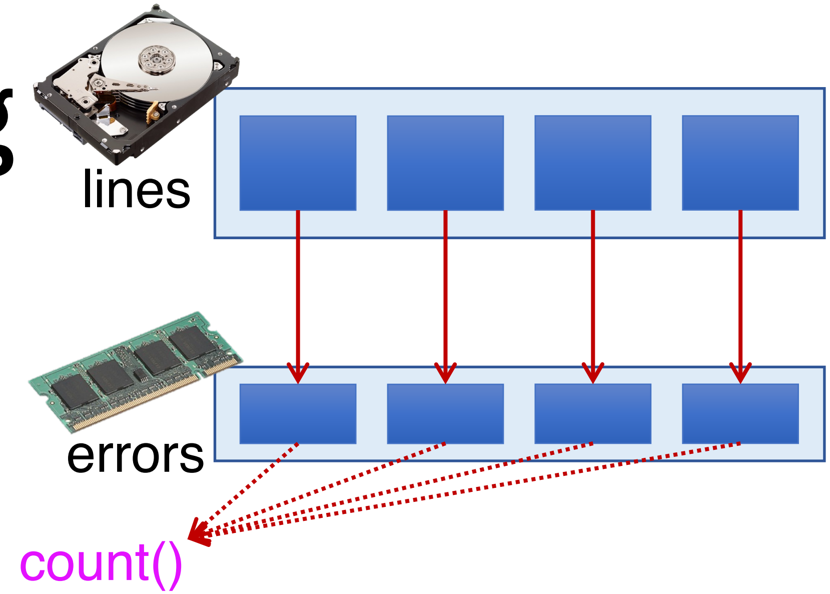

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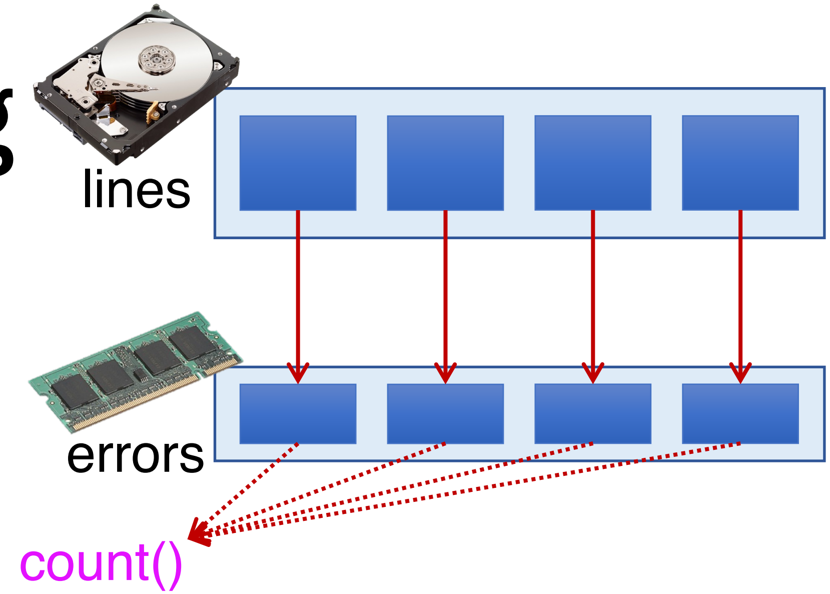

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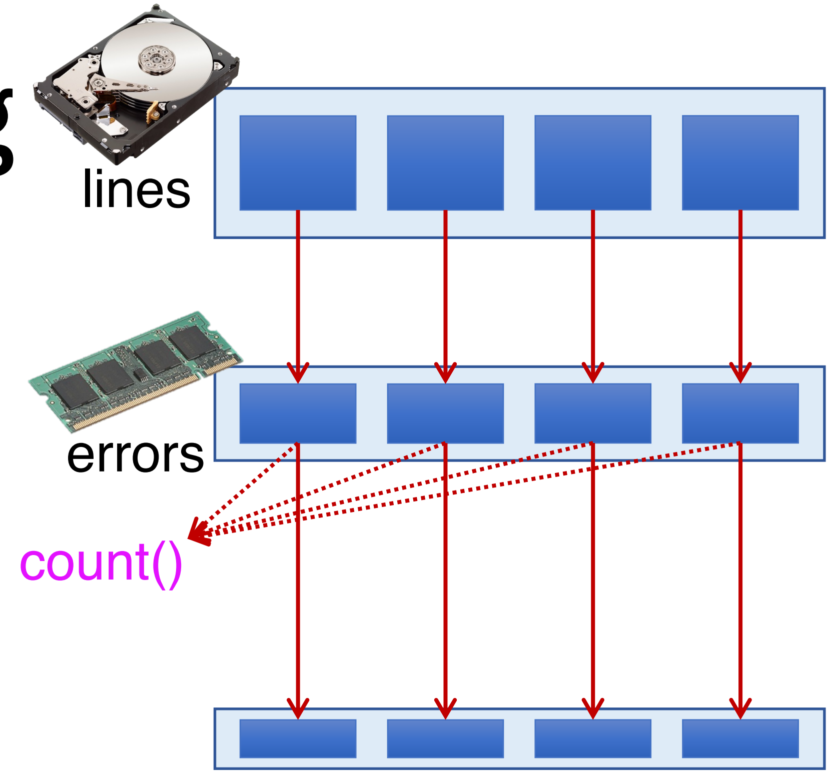

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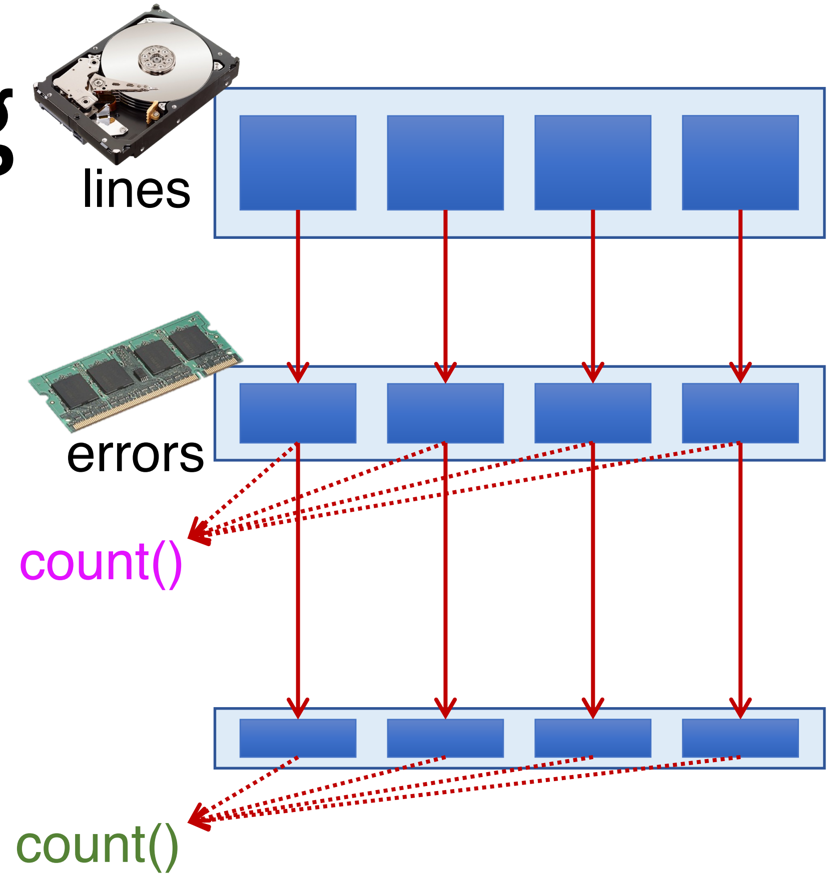

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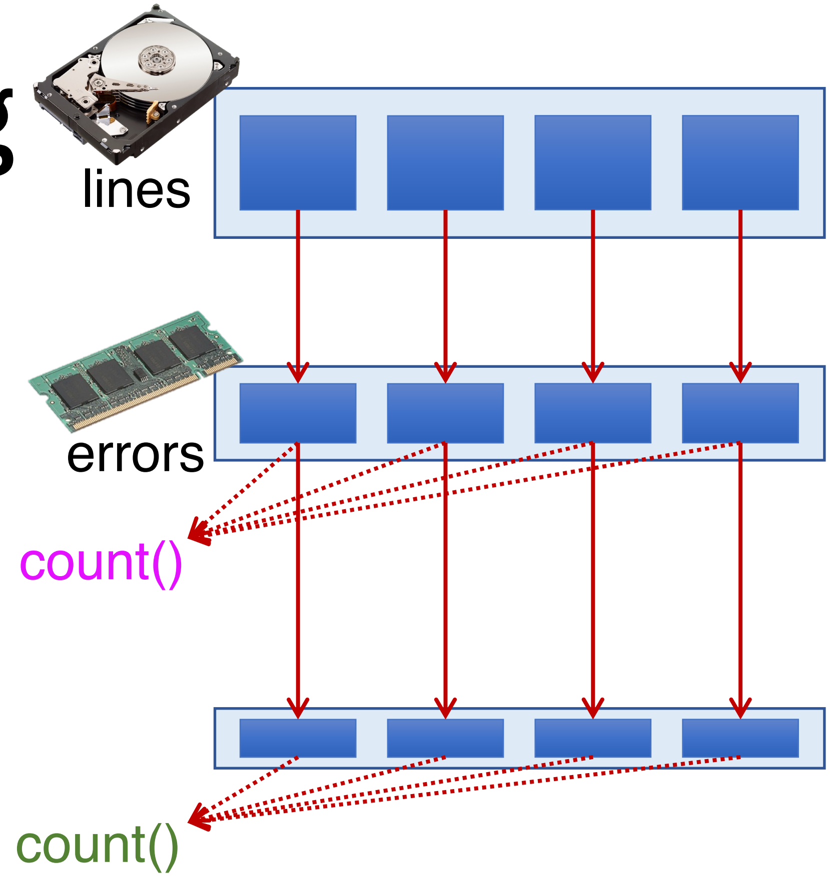

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errors.filter(
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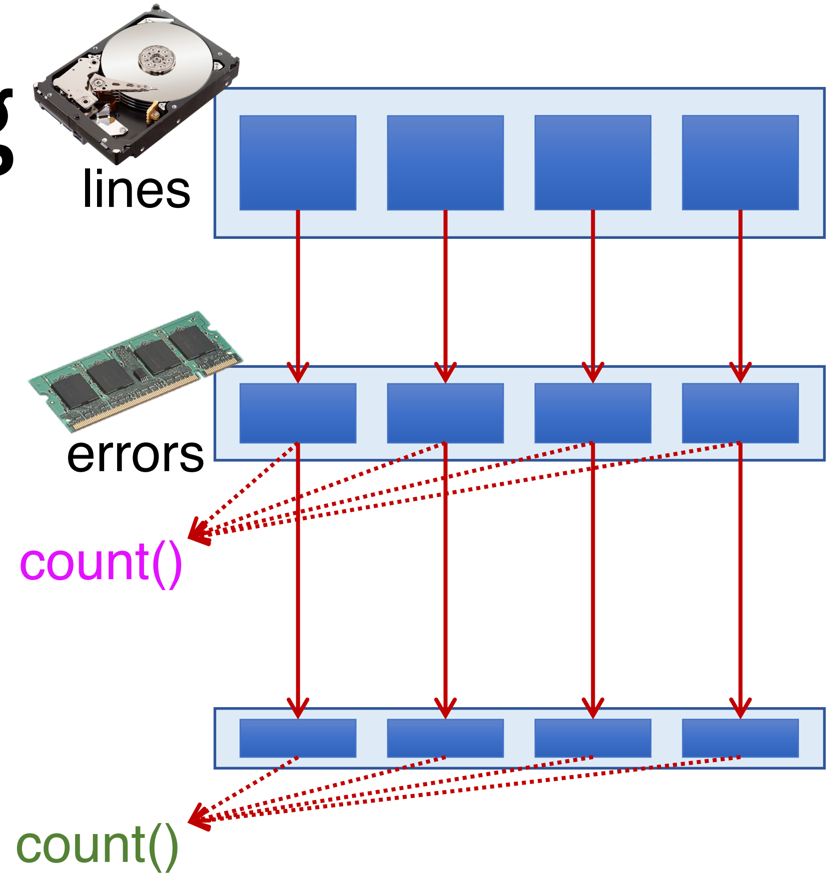

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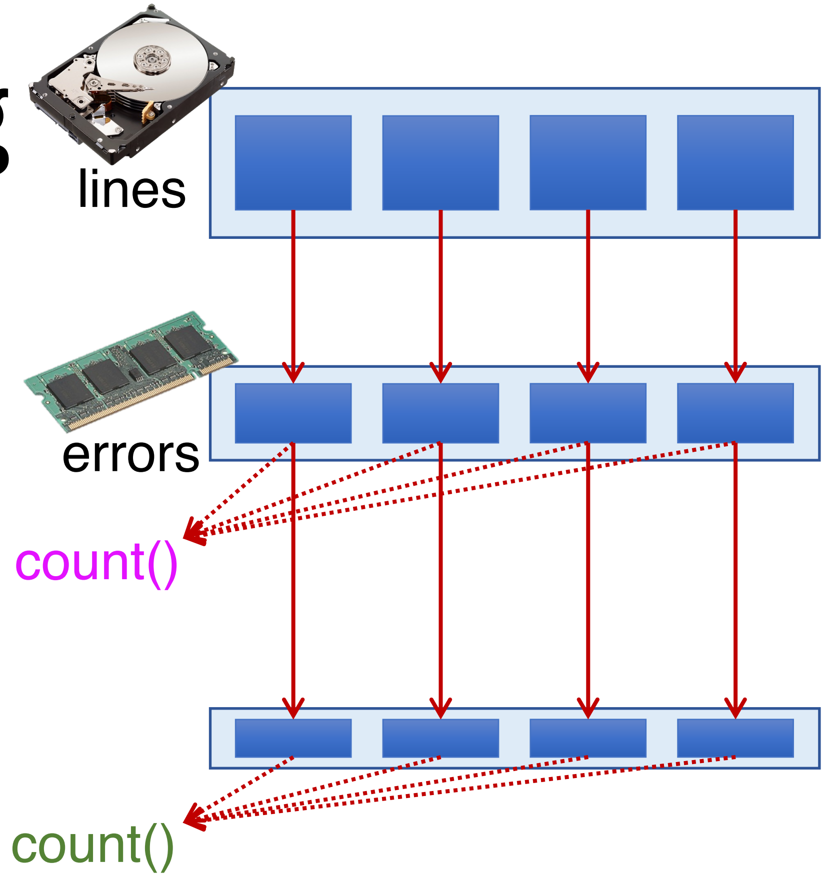

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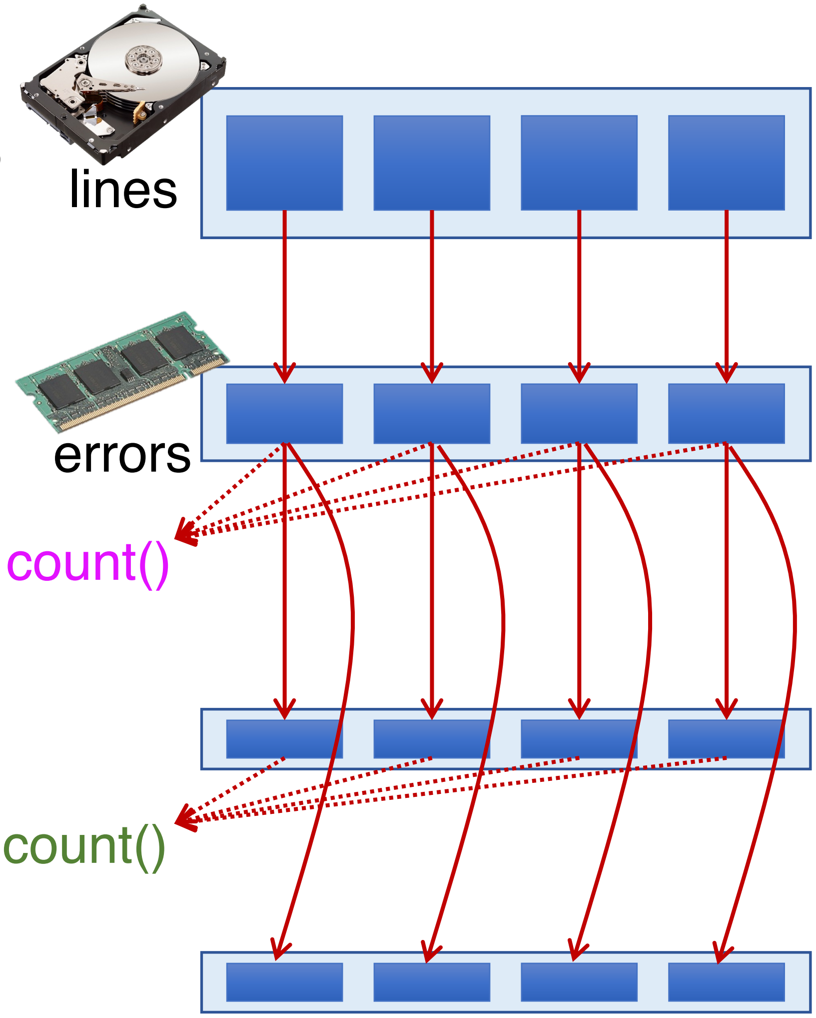

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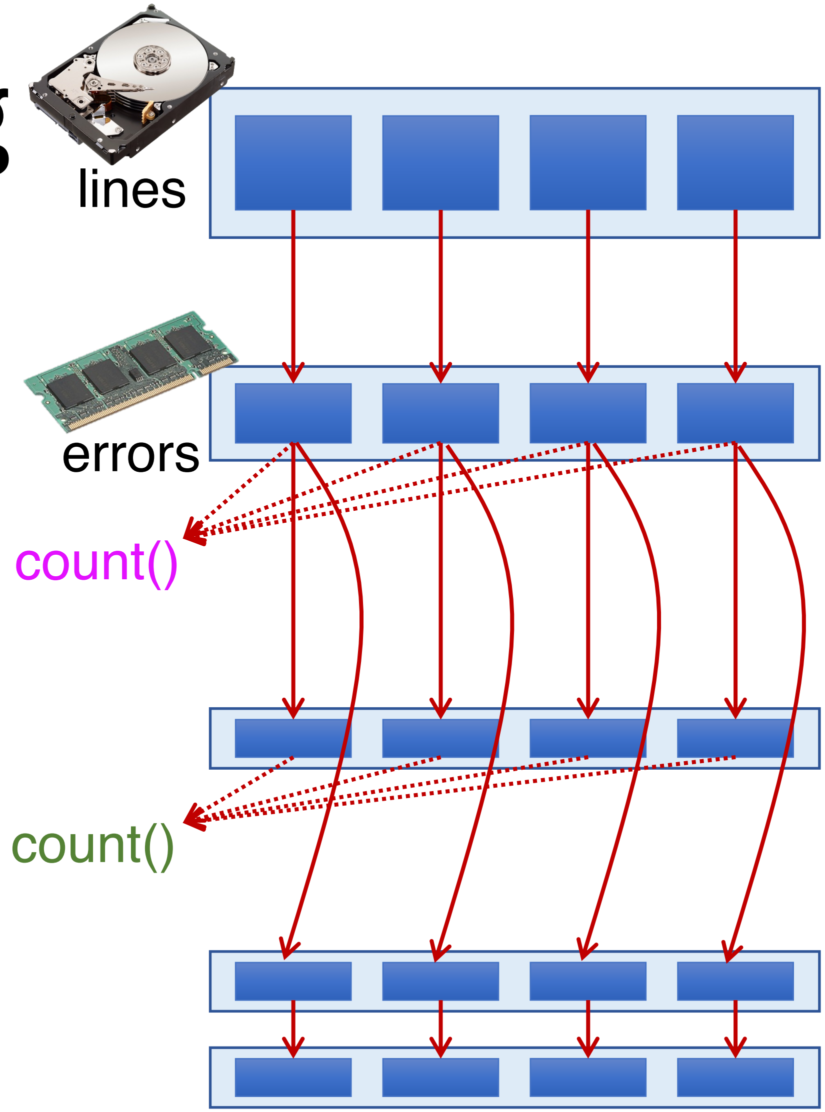

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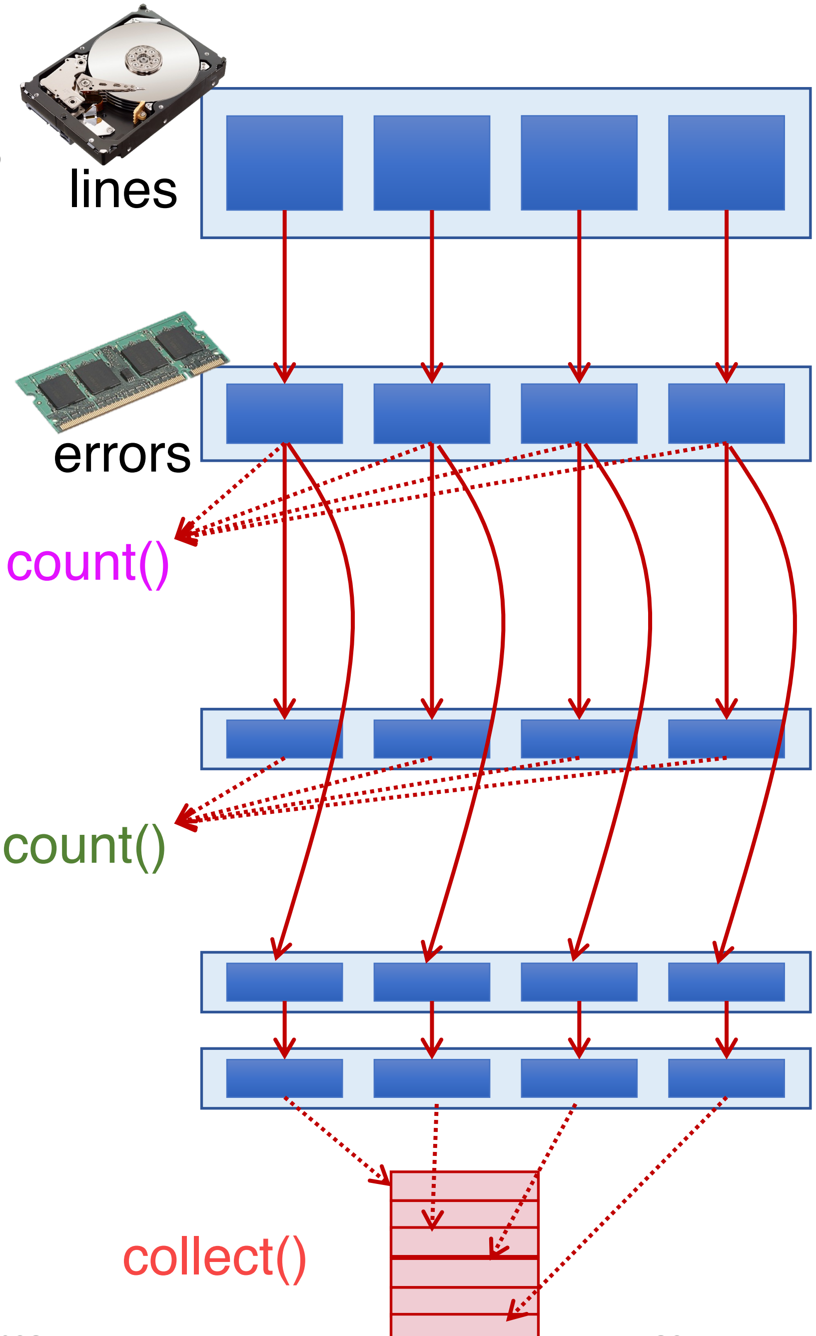


persist()

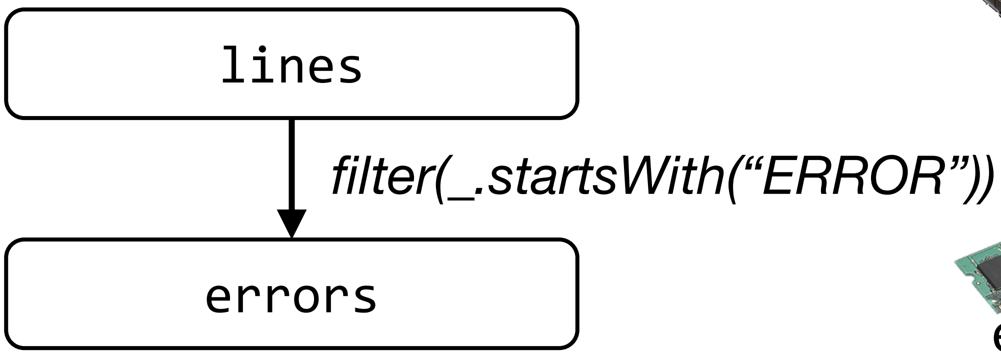
- Not an action nor a transformation
- A scheduler hint

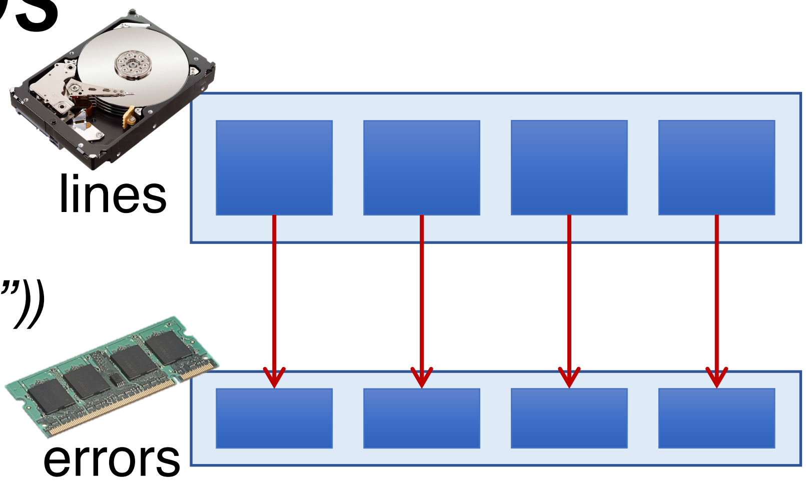
- Tells which RDDs the Spark schedule should materialize and whether in memory or storage
- Gives the user control over reuse/recompute/recovery tradeoffs

Lineage graph of RDDs

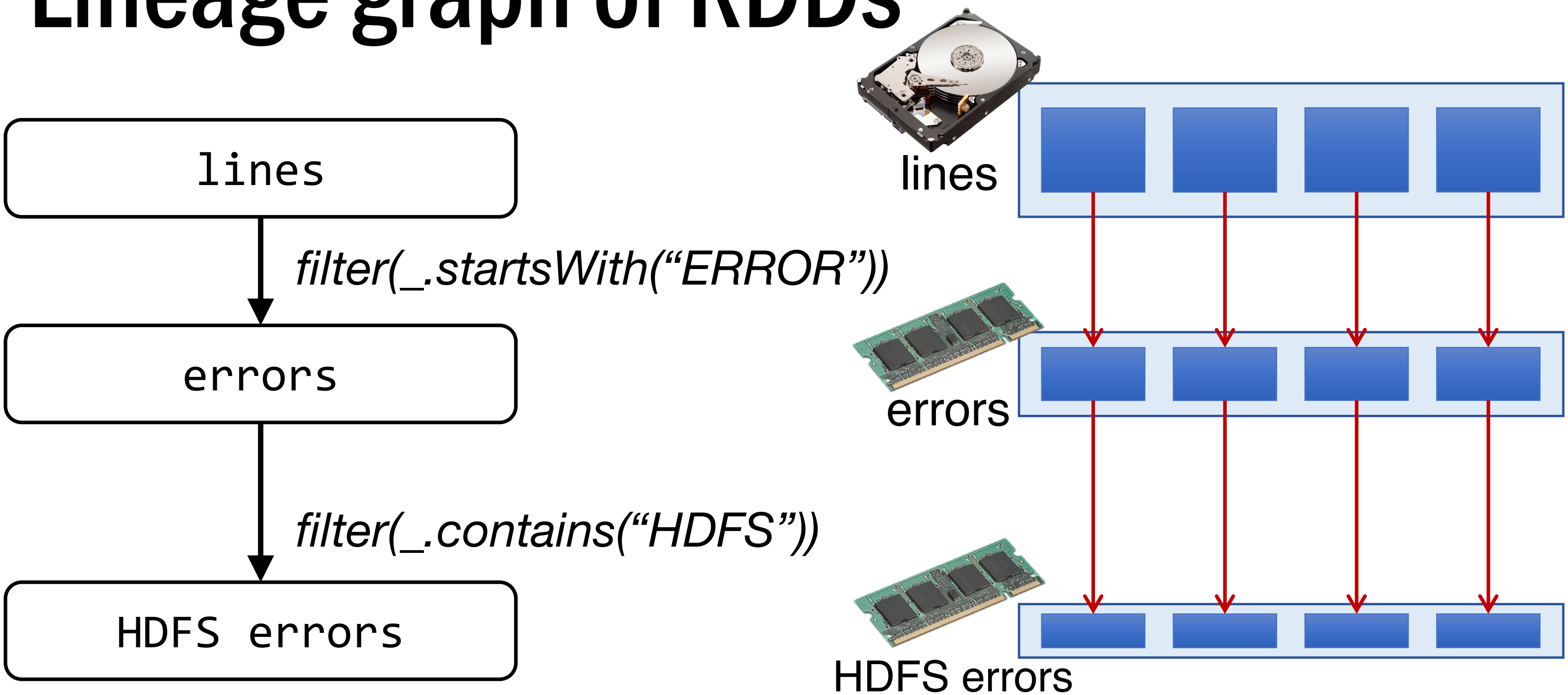
`lines`

Lineage graph of RDDs





Lineage graph of RDDs



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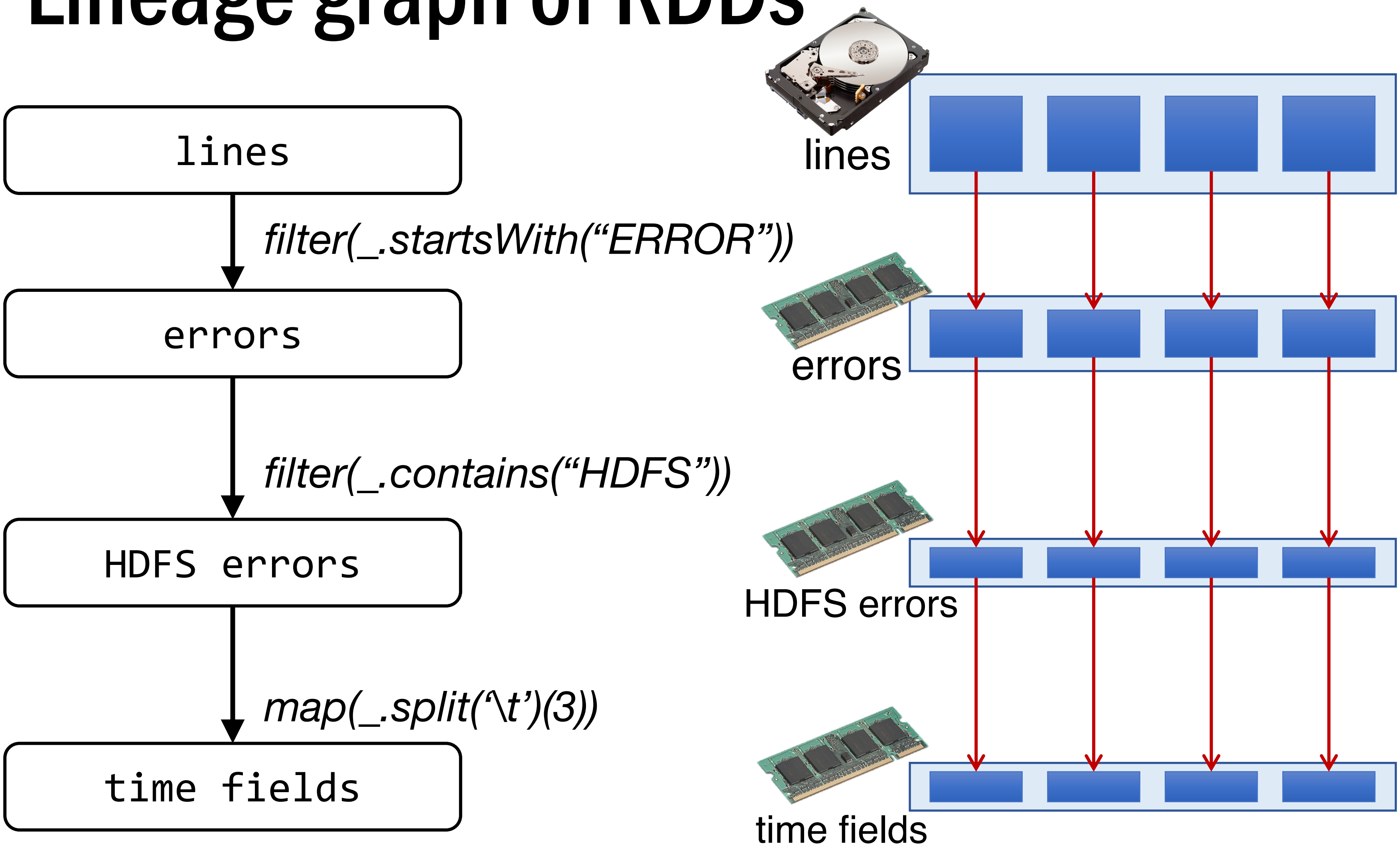
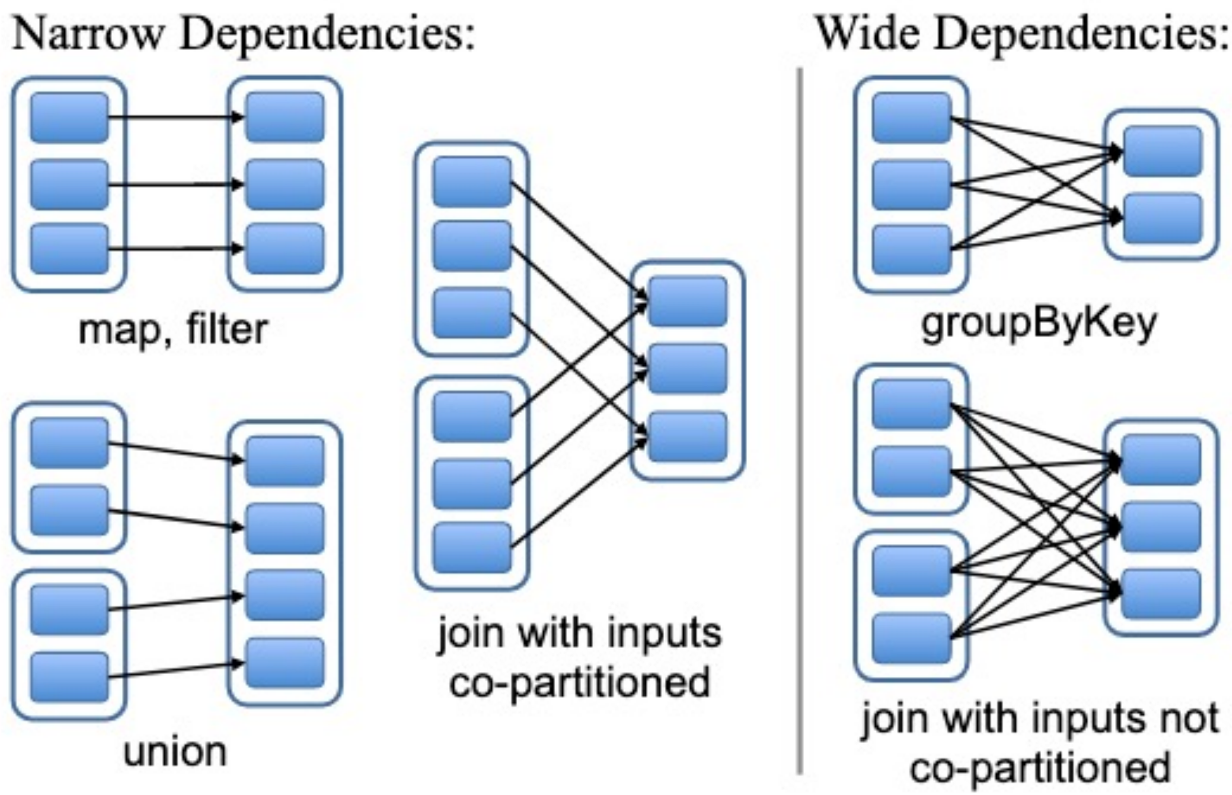


Narrow & wide dependencies



Narrow: each parent partition used by at most one child partition (can partition on one machine)

Wide: multiple child partitions depend on one parent partition

Must stall for all parent data, loss of child requires whole parent RDD (not just a small # of partitions)

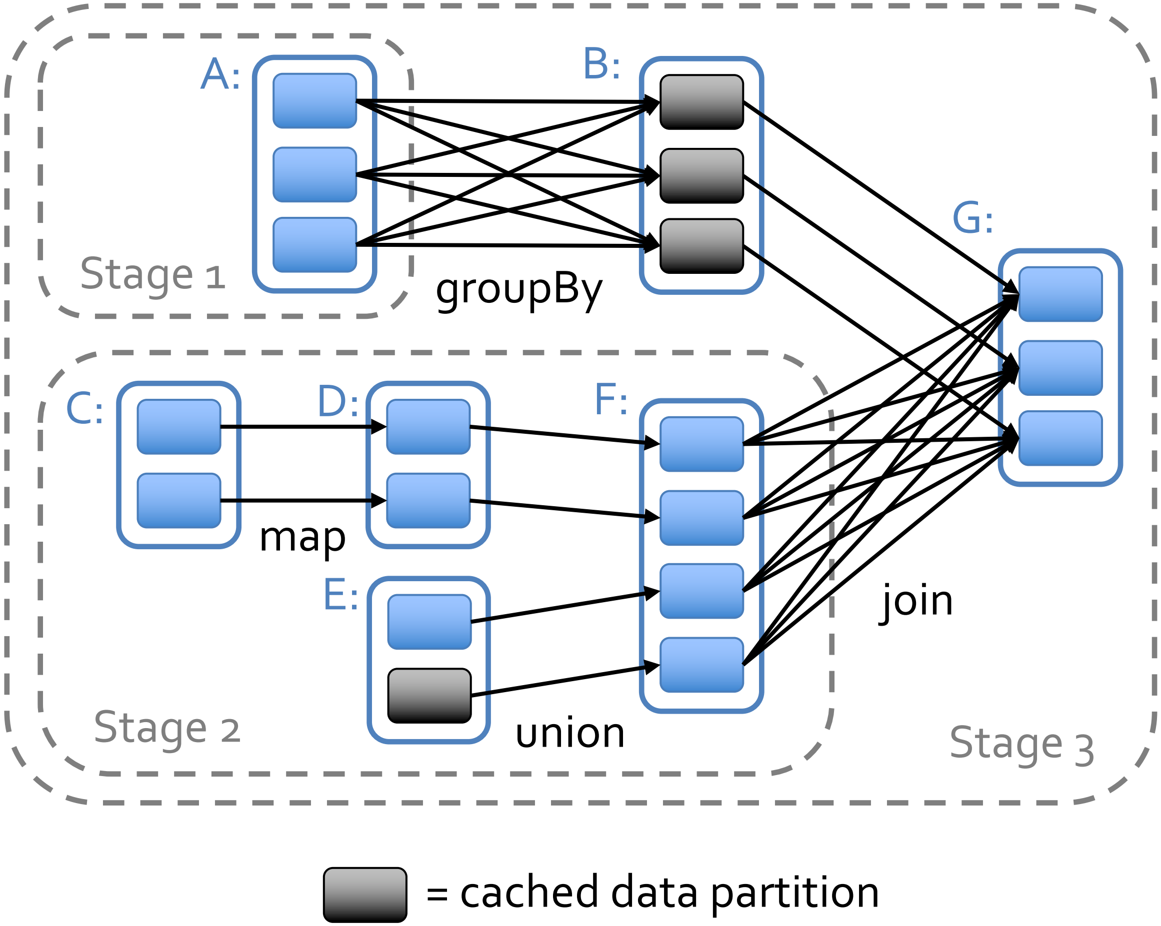
Task scheduler

DAGs (directed acyclic graphs)

Pipelines functions within a stage

Locality & data reuse aware

Partitioning-aware to avoid shuffles

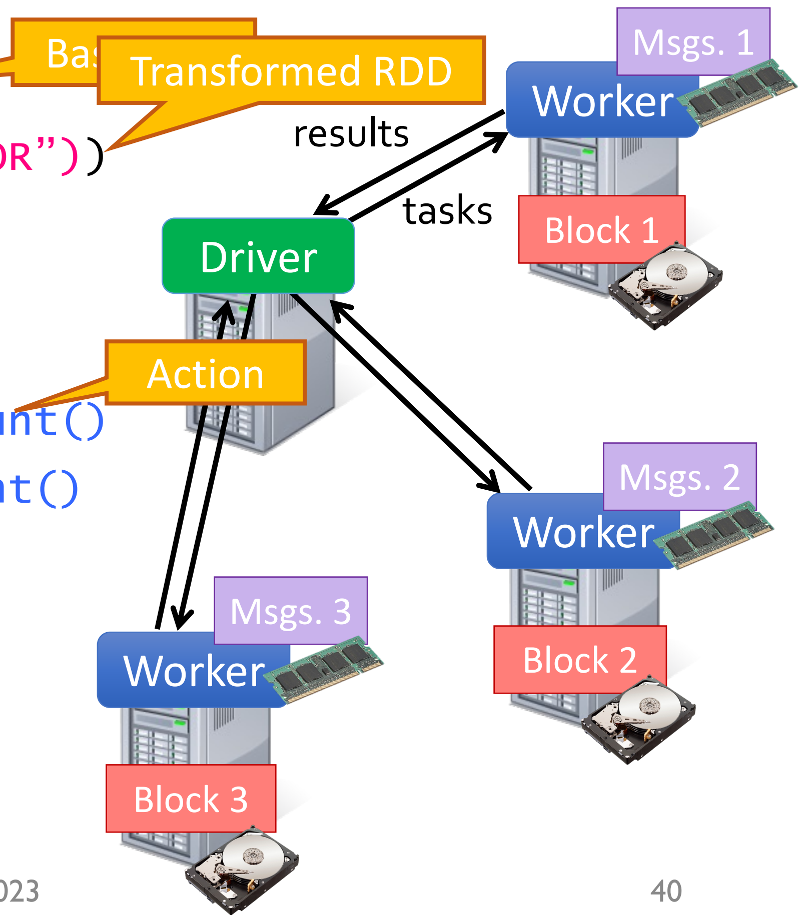


Interactive debugging (control and data flow)

Load error messages from a log into memory, then interactively search for various patterns

```
lines = spark.textFile("hdfs://...")
errors = lines.filter(_.startsWith("ERROR"))
messages = errors.map(_.split('\t')(2))
messages.persist()

messages.filter(_.contains("MySQL")).count()
messages.filter(_.contains("HDFS")).count()
```



Result: scaled to 1 TB data in 5-7 sec
(vs 170 sec for on-disk data)

Fault recovery

- RDDs track the graph of transformations that built them (their *lineage*) to rebuild lost data

E.g.:

```
messages = textFile(...).filter(_.contains("error"))
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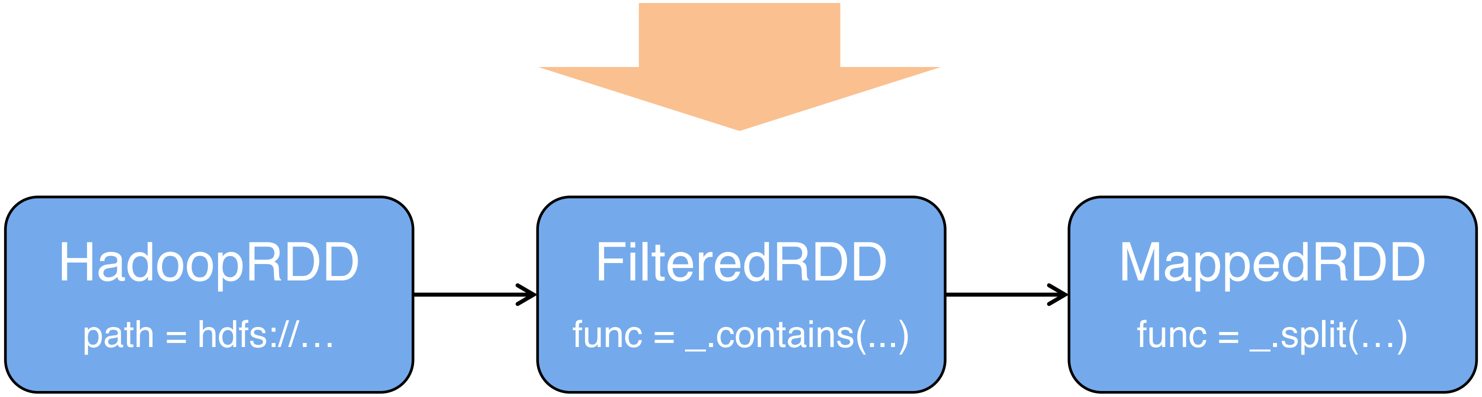

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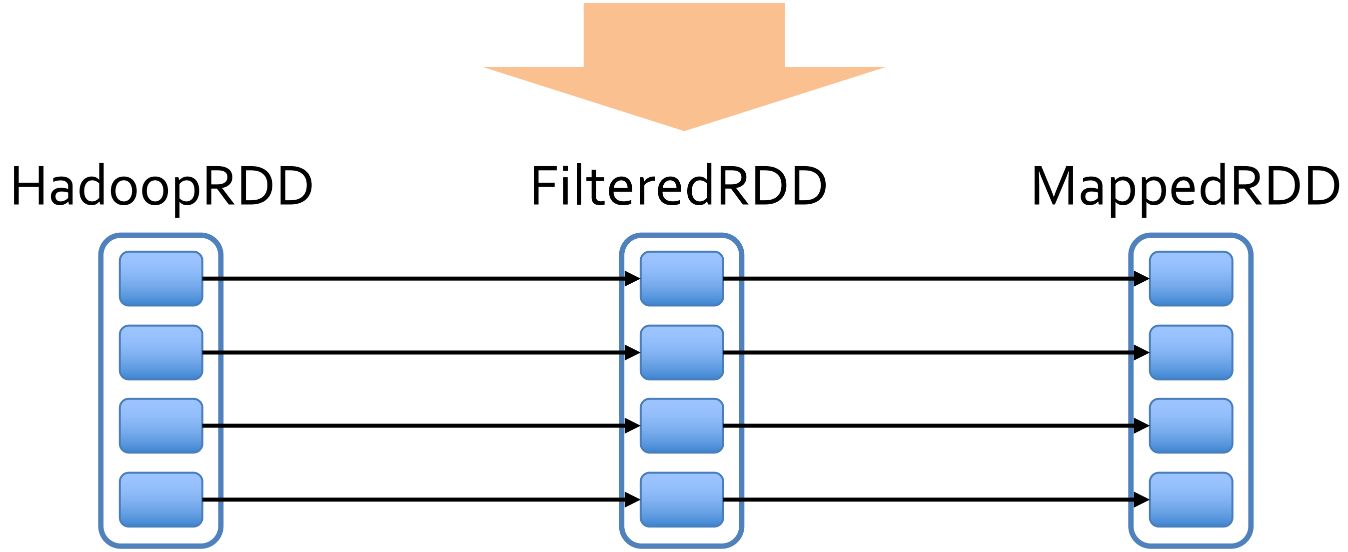

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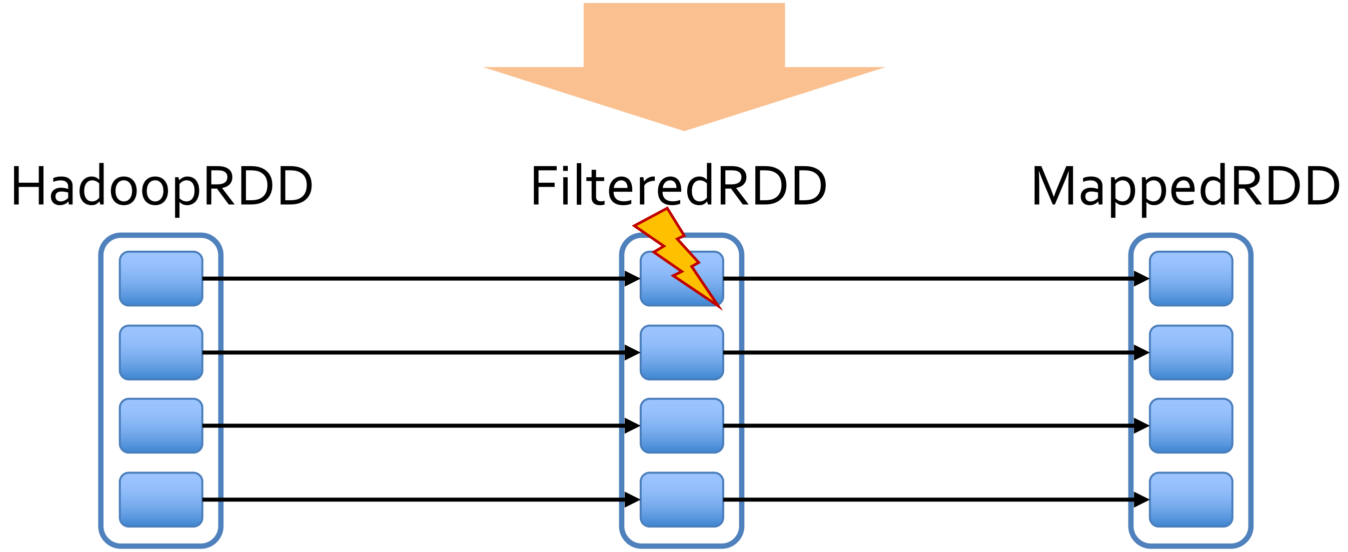
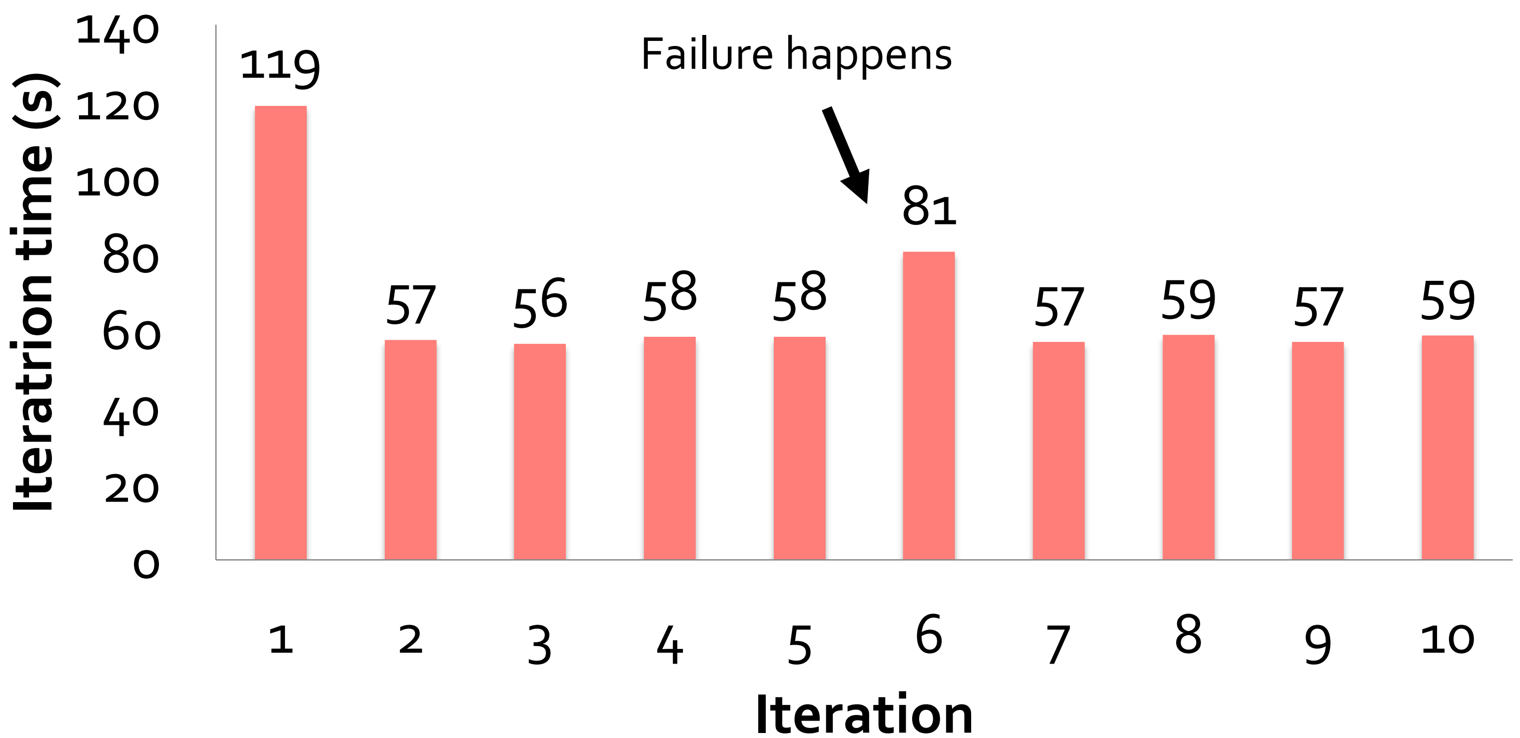


Fault recovery results



Example: PageRank

1. Start each page with a rank of 1
2. On each iteration, update each page's rank to

$$\Sigma_{i \in neighbors}\ rank_i\ /\ |neighbors_i|$$

```
links = // RDD of (url, neighbors) pairs
ranks = // RDD of (url, rank) pairs

for (i <- 1 to ITERATIONS) {
  ranks = links.join(ranks).flatMap {
    (url, (links, rank)) =>
      links.map(dest => (dest, rank/links.size))
  }.reduceByKey(_ + _)
}
```

Example: PageRank

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RDD[(URL, Seq[URL])]

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RDD[(URL, Rank)]

RDD[(URL, (Seq[URL],Rank))]

For each neighbor in links emits (URL, RankContrib)

Reduce to RDD[(URL, Rank)]

Join ($\bowtie$)

Alice	5
Bob	6
Claire	4

$\bowtie$

Alice	F
Bob	M
Claire	F

=

Alice	5	F
Bob	6	M
Claire	4	F

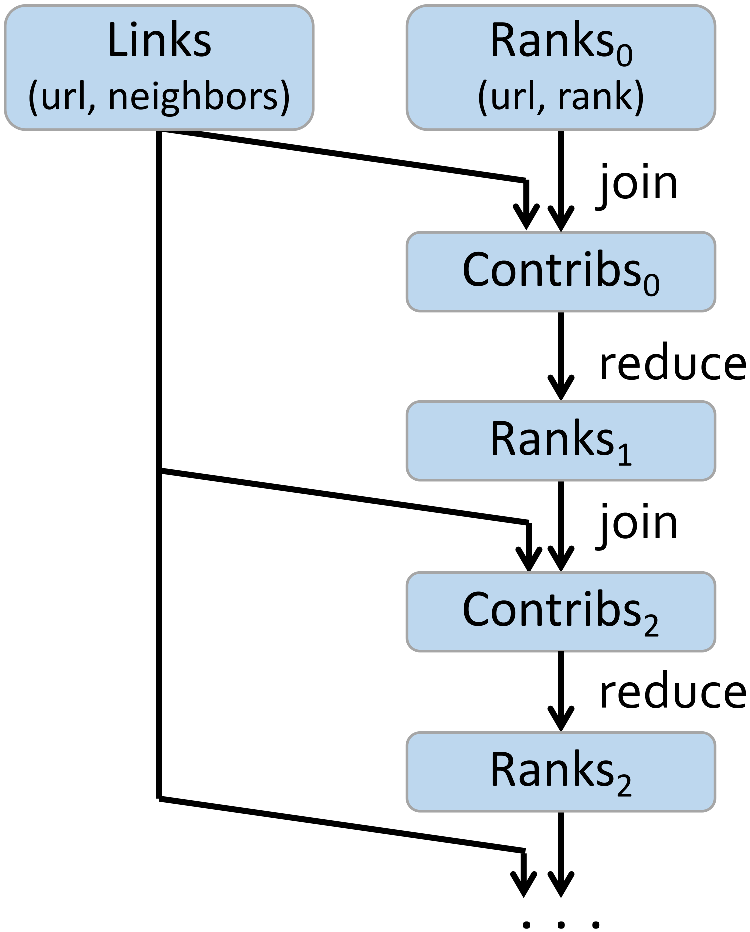
A	5
A	2
A	3
B	4
B	1
C	6
C	8

$\bowtie$

C	5
B	2
A	3
B	4
A	1
B	6
C	8

If partitioning doesn't match, then need to reshuffle to match pairs. Same problem in `reduce()` for MapReduce.

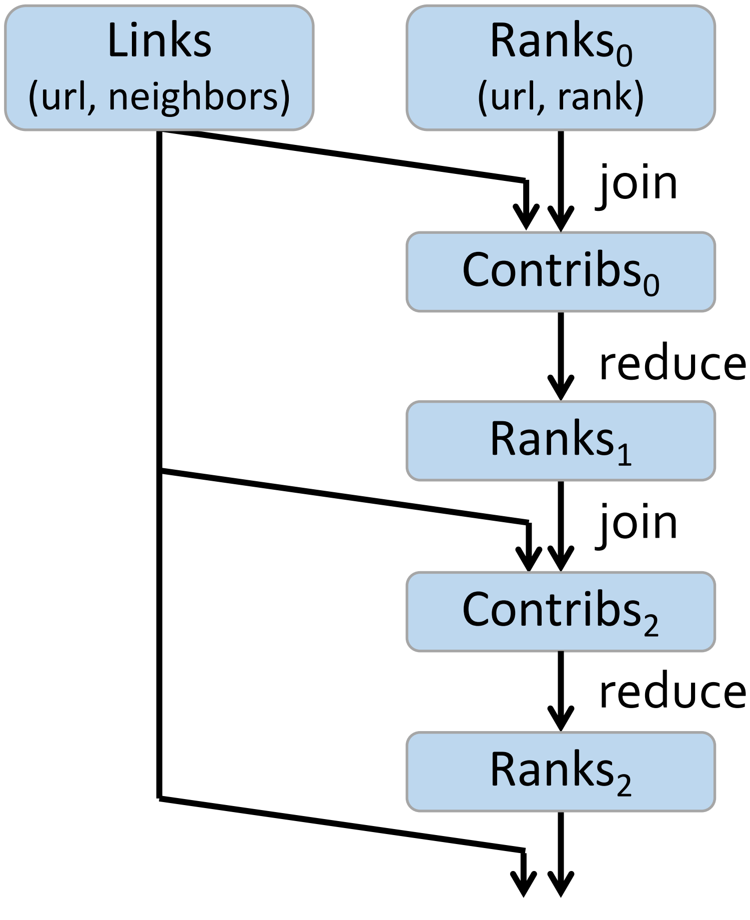
Optimizing placement



- `links` & `ranks` repeatedly joined
- Can *co-partition* them (e.g. hash both on URL) to avoid shuffles
- Can also use app knowledge, e.g., hash on DNS name

- ```
  links = links.partitionBy(
               new URLPartitioner())
  ```
  ```

Optimizing placement



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```

Q: Where might we have placed **`persist()`**?

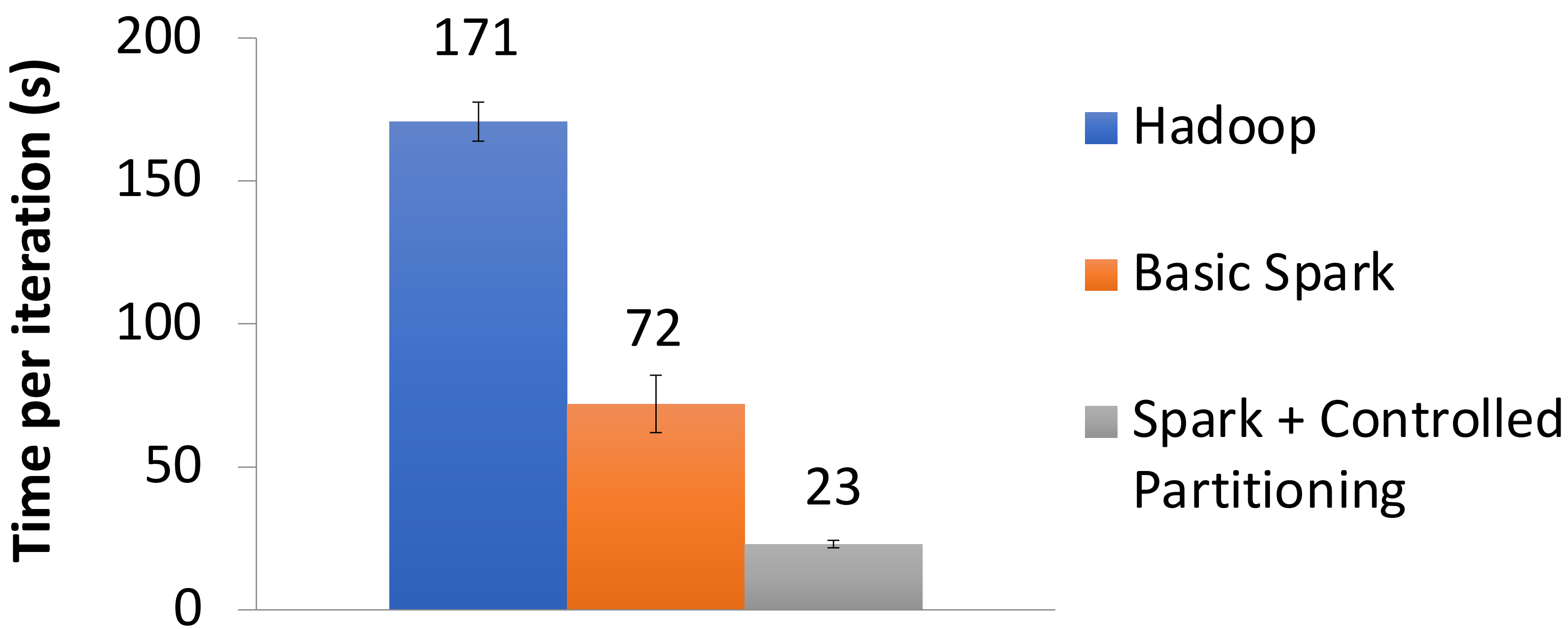
Co-partitioning example

Co-partitioning can avoid shuffle on join

But, fundamentally a shuffle on `reduceByKey`

Optimization: custom partitioner on domain

PageRank performance



* Figure 10a: 30 machines on 54 GB of Wikipedia data computing PageRank