Spark Internals and Performance

DS 5110/CS 5501: Big Data Systems Spring 2024 Lecture 5b

Yue Cheng



Some material taken/derived from:

• Wisconsin CS 320 by Tyler Caraza-Harter.

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Learning objectives

- Know different storage and caching levels of Spark RDD
- Understand basic hash partitioning and join operations
- Describe the Spark implementation of the PageRank graph mining algorithm and roles various Spark optimizations play

Collecting data

Collecting data (OK)



Collecting data (OK)



Collecting data (OK)



Collecting data (bad)



Collecting data (bad)





df refers to CSV file
df2 = df.where(???)

Scenario: want to do lots of computations on df2 **Goal:** avoid repeatedly reading HDFS and filtering df







Documentation snippet: <u>https://spark.apache.org/docs/2.2.2/tuning.html#memory-tuning</u> By default, Java objects are fast to access, but <u>can easily consume a factor of **2-5x** more space</u> than the "raw" data inside their fields. This is due to several reasons:

- Each distinct Java object has an "object header", which is about 16 bytes and contains information such as a pointer to its class. For an object with very little data in it (say one Int field), this can be bigger than the data.
- Java Strings have about 40 bytes of overhead over the raw string data (since they store it in an array of Chars and keep extra data such as the length), and store each character as two bytes due to String's internal usage of UTF-16 encoding. Thus a 10-character string can easily consume 60 bytes.
- Common collection classes, such as HashMap and LinkedList, use linked data structures, where there is a "wrapper" object for each entry (e.g. Map.Entry). This object not only has a header, but also pointers (typically 8 bytes each) to the next object in the list.
- Collections of primitive types often store them as "boxed" objects such as java.lang.Integer.



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Documentation snippet: https://spark.apache.org/docs/2.2.2/tuning.html#serialized-rdd-storage When your objects are still too large to efficiently store despite this tuning, a much simpler way to reduce memory usage is to store them in *serialized* form, using the serialized StorageLevels in the RDD persistence API, such as **MEMORY_ONLY_SER**. Spark will then store each RDD partition as one large byte array. The only downside of storing data in serialized form is slower access times, due to having to deserialize each object on the fly.





Join and partitioning

Join and partitioning (best case)

Computer 1	Alice	5		Alice	F		Alice	5	F
	Bob	6	\bowtie	Bob	Μ	=	Bob	6	Μ
Computer 2	Claire	4		Claire	F		Claire	4	F

Join and partitioning (worst case)

Computer 1	Alice	5		Alice	F			Alice	5	F		
	Bob	6	\bowtie	Bob	Μ		=	Bob	6	Μ		
Computer 2	Claire	4		Claire	F			Claire	4	F		
	А	5		С		5						
	А	2		В		2						
Computer 1	А	3		A		3		If partitioning does match, then need to shuffle (via networ to match pairs				
Computer 2	В	4	Å	₫В		4						
	В	1	K	A		1		to match pairs.				
	С	6		В		6						
	С	8		► C		8						

Join and partitioning (optimization)

Computer 1	Alice	5		Alice	F			Alice	5	F	
	Bob	6	\bowtie	Bob	Μ		=	Bob	6	Μ	
Computer 2	Claire	4		Claire	F			Claire	4	F	
	A	5									
	А	2		A		3		par ⁻	artitionBy() is		
Computer 1	А	3		A		1		pair RDDs. It is us			
Computer 2	В	4		В		2		to p base	ed on keys, by ault using a hash		
	В	1		B		4		defa			
	С	6		В		6		partitioner.			
	С	8		C		5					
				С		8					

Example: PageRank

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1. Start each page with a rank of 1 2. On each iteration, update each dest page's rank to $\Sigma_{i \in neighbors}$ rank_{neighbor_i} / |neighbors_i|

```
links = // RDD of (url, neighbors) pairs
ranks = // RDD of (url, rank) pairs
for (i <- 1 to ITERATIONS) {
  ranks = links.join(ranks).flatMap {
    (url, (links, rank)) =>
    links.map(dest => (dest, rank/links.size))
  }.reduceByKey(_ + _)
}
```

Example: PageRank

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Demo ...

Optimizing placement



- links & ranks repeated joined
- Can co-partition them (e.g., hash both on source URLs) to avoid shuffles

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Q1: Should we apply .persist(DISK_ONLY) to links or ranks?

Q2: Where might we have placed .persist(DISK_ONLY)?

Discussion: Spark perf (paper)

