

Memory Management: Page Replacement Policies: Miscellaneous Topics

CS 571: Operating Systems (Spring 2020)

Lecture 8c

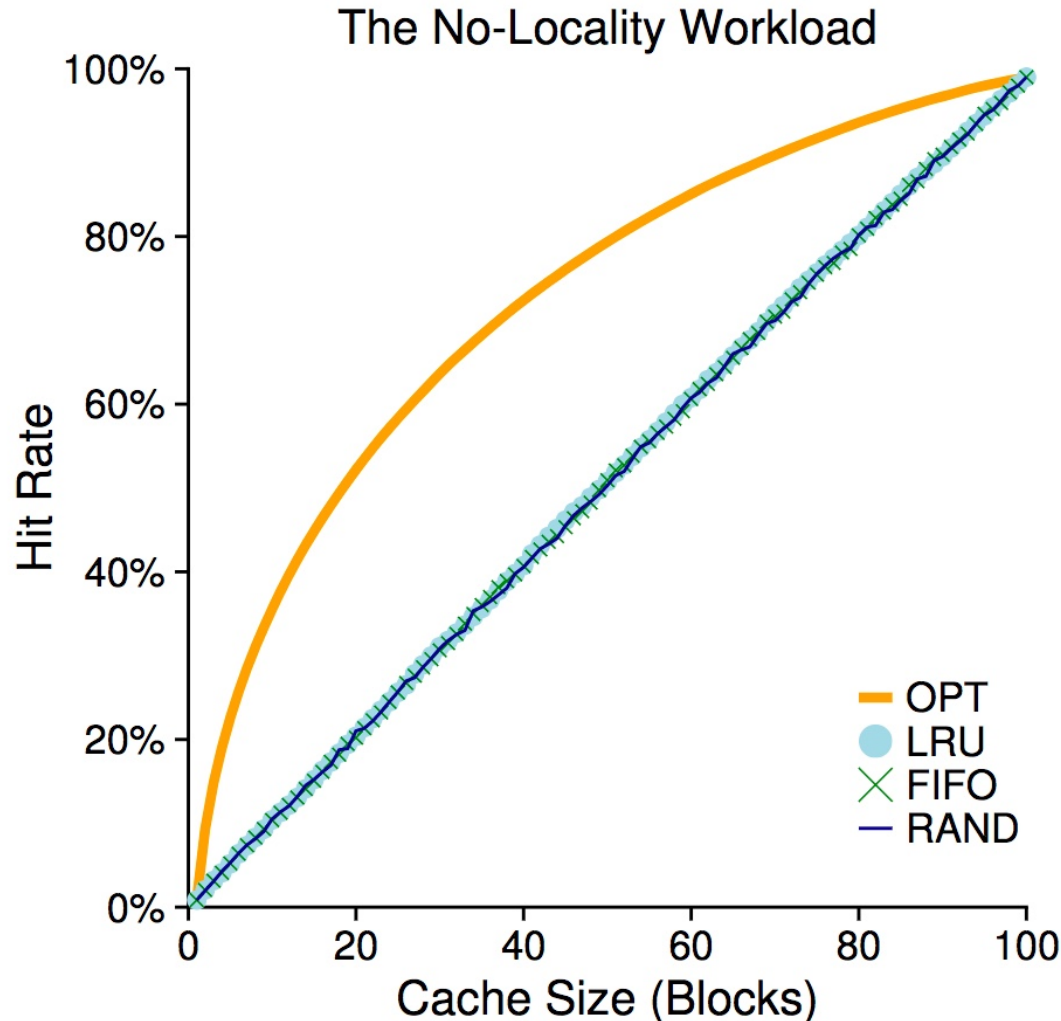
Yue Cheng

Page Replacement Workload Examples

Workload Examples

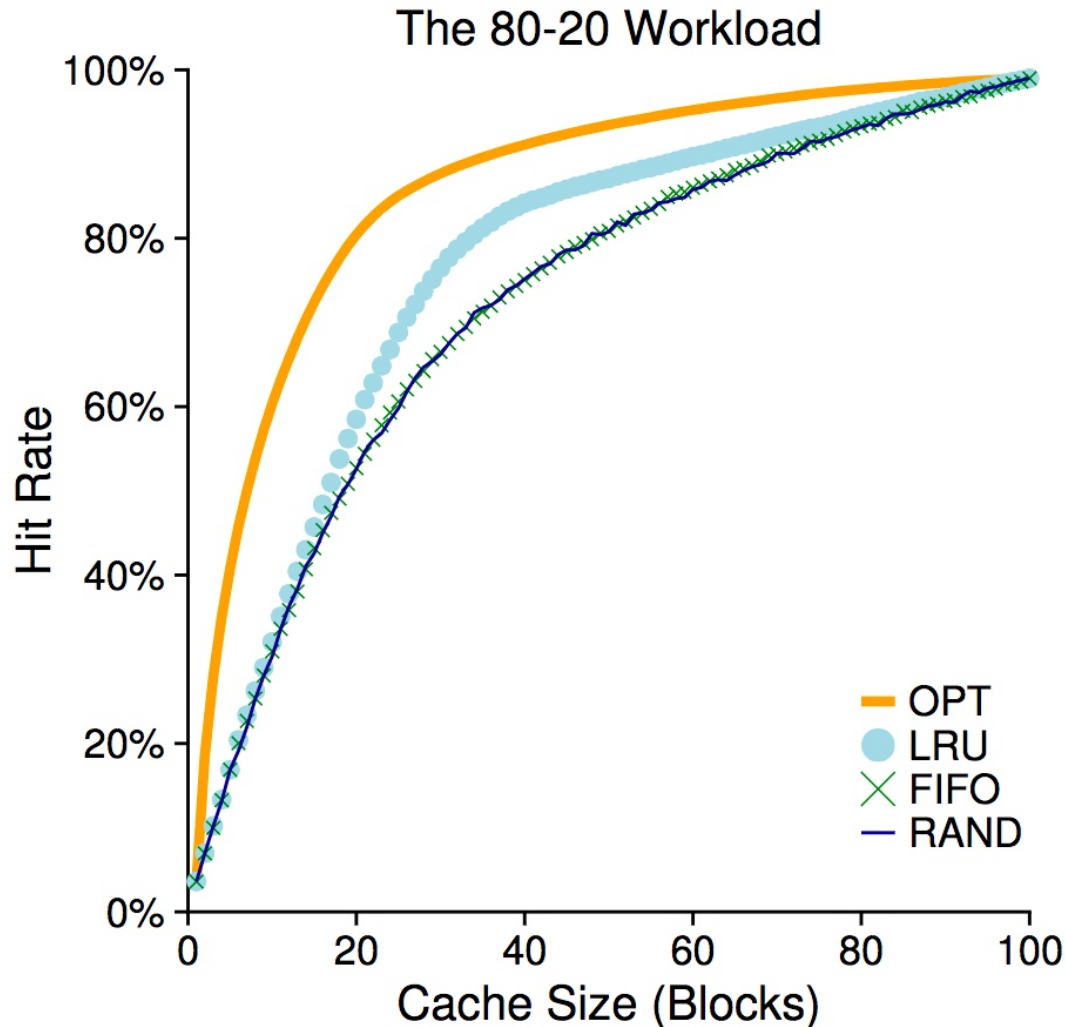
- A simple workload
 - Workload consists of a working set of 100 pages
 - Workload issues 10,000 access requests
- Four replacement policies
 - OPT: The optimal
 - LRU: Least-recently used
 - FIFO: First-in first-out
 - RAND: Random

The No-Locality Workload



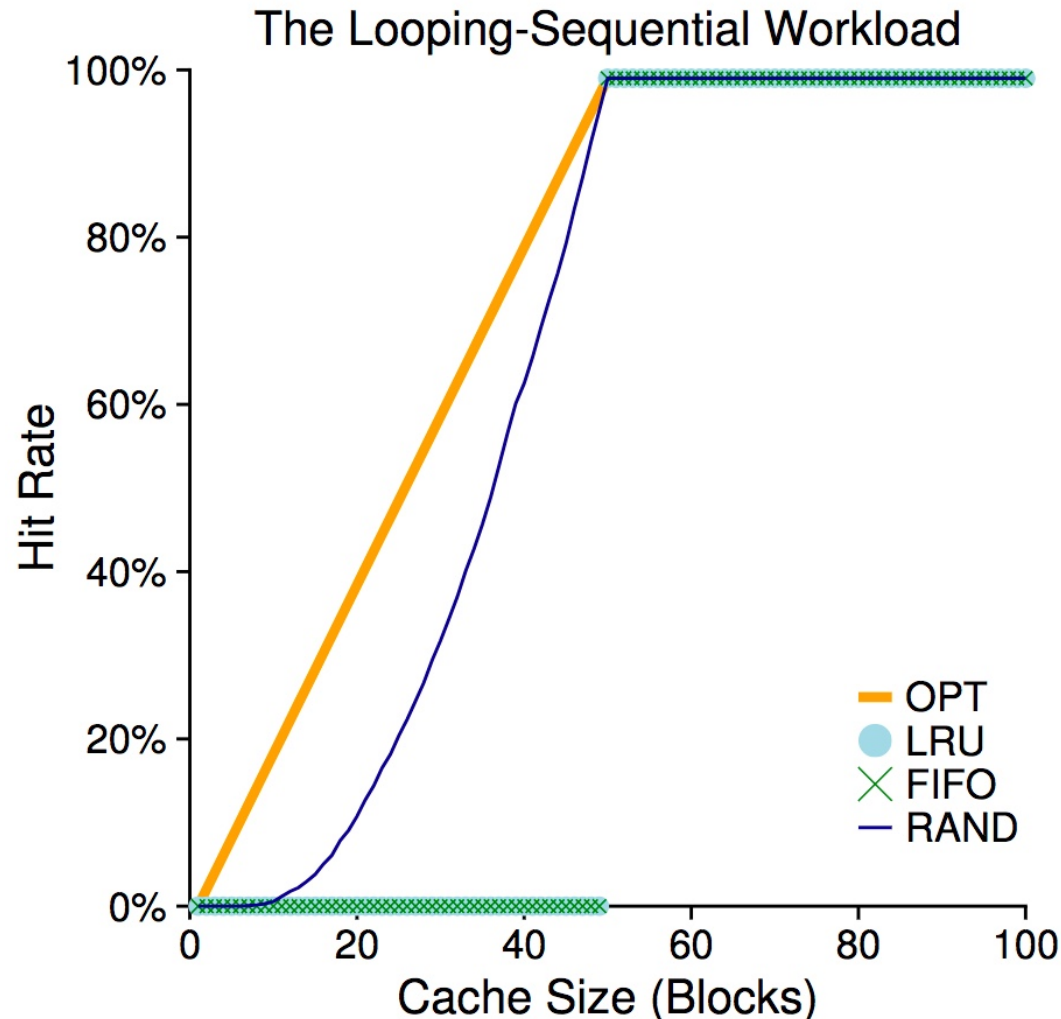
Each reference is to a random page within the set of accessed pages

The 80-20 Workload



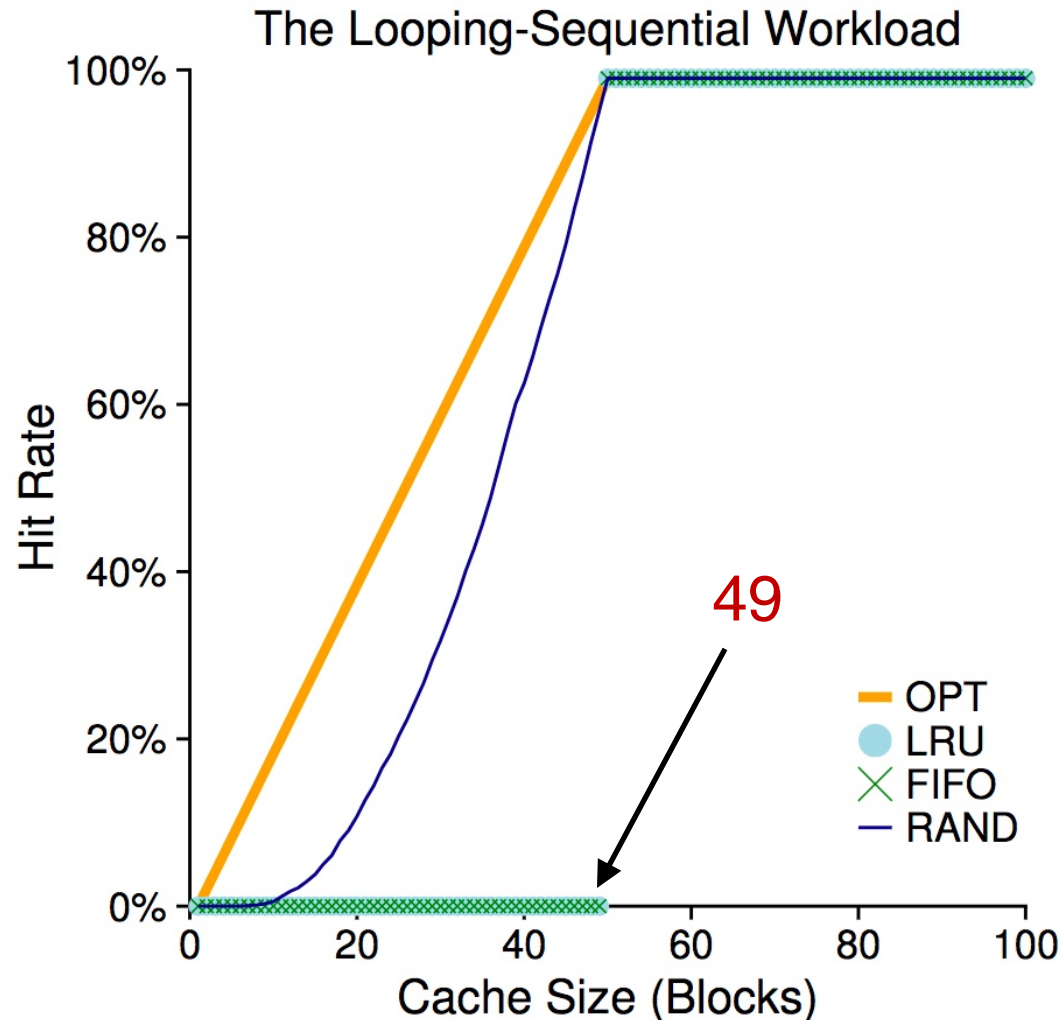
80-20: 80% of the refs are made to 20% of the pages (“**hot**” pages)

The Looping-Sequential Workload



Loop first 50 pages starting from 0 to 49 for a total of 10,000 accesses

The Looping-Sequential Workload



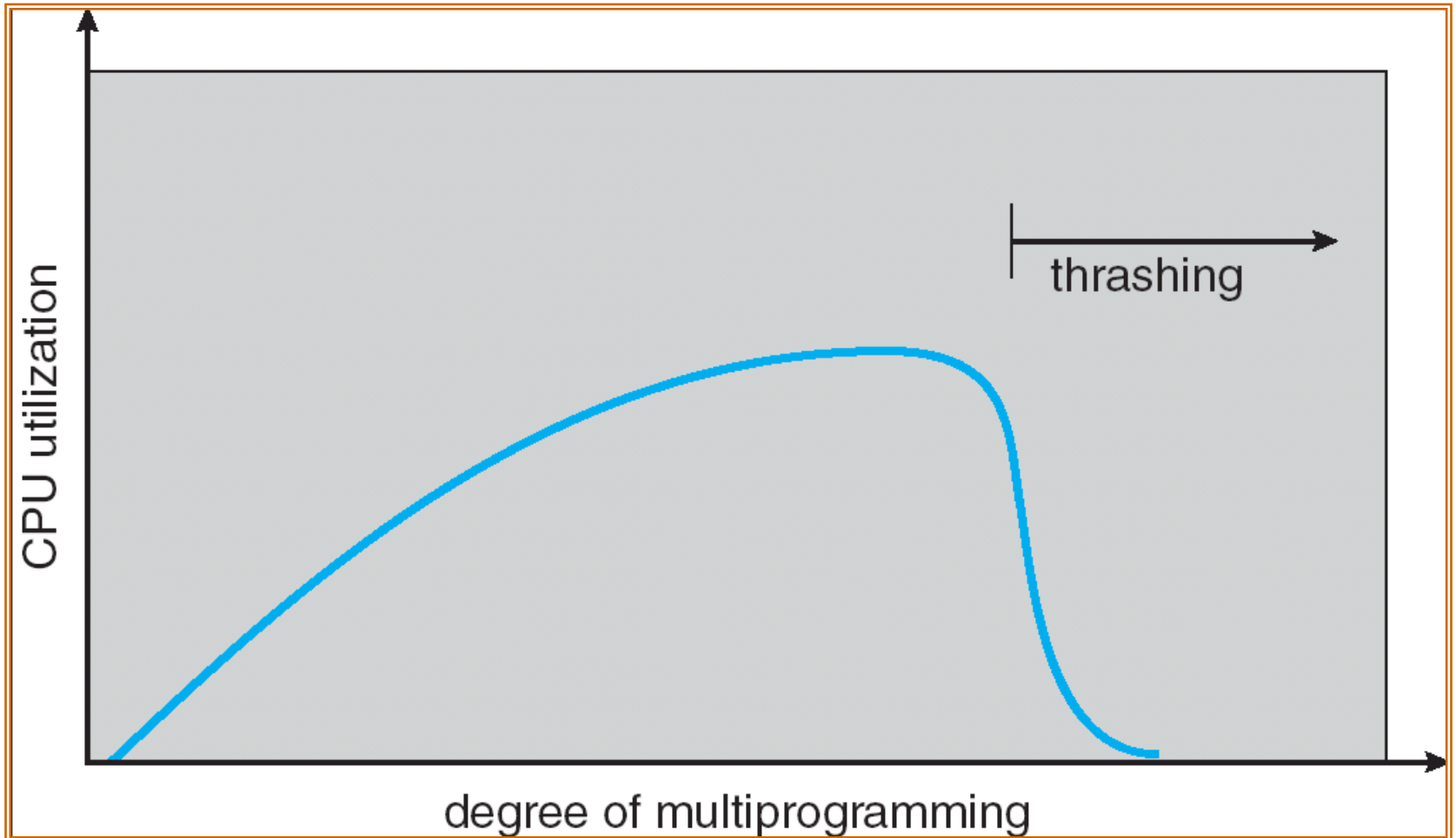
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Thrashing

Thrashing

- High-paging activity: The system is spending more time paging than executing
- How can this happen?
 - OS observes low CPU utilization and increases the degree of multiprogramming
 - Global page-replacement algorithm is used, it takes away frames belonging to other processes
 - But these processes need those pages, they also cause page faults
 - Many processes join the waiting queue for the paging device, CPU utilization further decreases
 - OS introduces new processes, further increasing the paging activity

CPU Utilization vs. the Degree of Multiprogramming



How to Avoid Thrashing?

- To avoid thrashing, earlier OS did admission control to only run a subset of processes
- Some current OS takes more draconian approach
 - E.g., some Linux runs an out-of-memory killer to choose a memory-intensive process and kill it

Review: Demand Paging

- Bring a page into memory **only when it is needed**
 - Less I/O needed
 - Less memory needed
 - Faster response
 - Support more processes/users
- Page is needed \Rightarrow use the reference to page
 - If not in memory \Rightarrow must bring from the disk
- Demand paging versus swapping
 - Fetching the page in only on demand vs. kicking out one victim then paging in one under mem pressure

Demand Paging and Thrashing

- Why does demand paging work?
Locality model
 - Process migrates from one locality to another
 - Localities may overlap
- **Why does thrashing occur?**
 Σ size of locality > total memory size
Or Σ working set size > total memory size
- Definition of **working set size** (WSS): number of unique items that are accessed

Impact of Program Structures on Memory Performance

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- Program 1

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for (j = 0; j < 128; j++)  
    for (i = 0; i < 128; i++)  
        data[i][j] = 0;
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$128 \times 128 = 16,384$ page faults

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- Program 2

```
for (i = 0; i < 128; i++)  
    for (j = 0; j < 128; j++)  
        data[i][j] = 0;
```

Only **128** page faults