# OS161 Project 2

System Calls and Process Scheduling

## Deliverables

- Answers to the code walk-through questions
- Design Document
- Implementations
  - System calls
    - getpid
    - fork
    - execv
    - waitpid
    - exit
  - Multi-level queue scheduler

## Deliverables (contd.)

#### Design document

- A high level description of how you are approaching the problem
- A detailed description of the implementation (e.g., new structures, why they were created, what they are encapsulating, what problems they solve)
- A discussion of the pros and cons of your approach
- Alternatives you considered and why you discarded them

## **Configure and Build Kernel**

Repeat the steps you used for the last project

• Just use ASST2 instead of ASST1

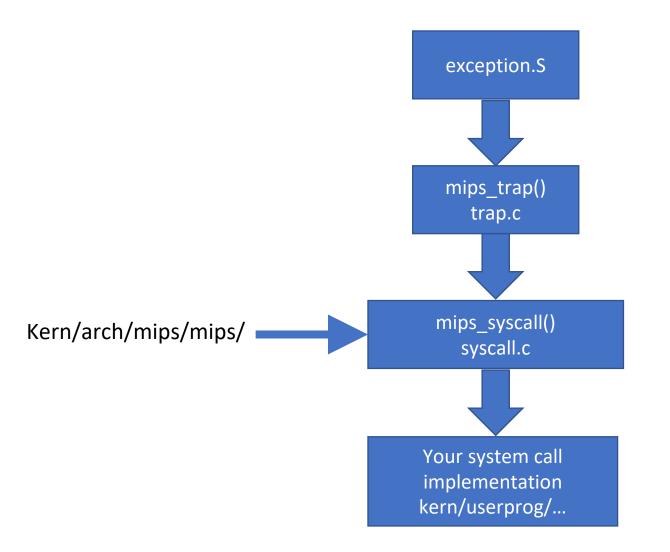
## Where to put your system call implementation?

- This time no skeleton code is given
- Create under kern/userprog
  - fork.c
  - execv.c
  - waitpid.c
  - getpid.c
  - exit.c
- Name your system calls sys\_{getpid|fork|execv|waitpid|exit}
- Add the new files to kern/conf/conf.kern
  - e.g., file userprog/getpid.c
  - The same way you have done hello.c in ASSTO.
  - Include your system call function declaration in kern/include/syscall.h

### **User-Level Interface**

- os161-1.11/include/unistd.h contains the user-level system call interfaces.
  - int execv(const char \*prog, char \*const \*args);
  - pid\_t fork(void);
  - int waitpid(pid\_t pid, int \*returncode, int flags);
  - int open(const char \*filename, int flags, ...);
  - int read(int filehandle, void \*buf, size\_t size);
  - int write(int filehandle, const void \*buf, size\_t size);
  - int close(int filehandle);
  - int reboot(int code);
  - int sync(void);

### How is it linked?



## For each system call

- mips\_syscall(struct trapframe \*tf) gets called
- The system call number is in tf->tf\_v0
- The arguments are in tf->tf\_a0, ..., tf->tf\_a3
- Case-switch statement calls the correct system call based on call number, and passes the arguments extracted from the trapframe
- ✓ Increment user-program counter before returning from system call Otherwise, it will restart the same system call tf->tf\_epc+=4
- ✓ If error

Store the error code in tf->tf\_v0 Set tf->tf\_a3 to 1

✓ If no error

Store the return value in tf->tf\_v0 Set tf->tf\_a3 to 0

#### Process structure

#### A common hack.

- Add the necessary fields to the thread structure and treat it as a process.
  - Pid
  - Exit status
  - Parent process
  - Etc.
- A process table
- A new pid needs to be generated for a new process
  - Need to reuse pid of processes that already exited

## Sys\_getpid

- Simplest one.
- Just return the pid of the executing process.
- getpid does not fail.

## Sys\_fork

- Duplicate the current process.
  - Child process will have unique process id.
- pid\_t sys\_fork(struct trapframe\*tf, pid\_t\*retval)
  - Child process returns 0.
  - Parent process return the pid of the child process.
- In case of an error
  - do not create a new child process but return -1.
- Most of the work is already done in thread.c (thread\_fork).
  Add the followings:
  - Create a pid when creating a new process. Add it to your process table.
  - Copy the trapframe.
  - Copy the address space.
  - Call thread\_fork()

## Sys\_fork

- Implement md\_forkentry
  - Parent's trapframe and address space are passed as arguments
  - Create new child trapframe by copying parent's
  - Get the assigned child pid from parent's trapframe tf\_v0 and assign it to the pid of the current process (since we are executing md\_forkentry, this is child)
  - Set the trapframe's tf\_v0 to 0.
  - Increment tf\_epc by 4.
  - Copy the passed address space to the current process address space and activate it.
  - Give the control back to the usermode.
    - Call mips\_usermode() and pass the new trapframe.

## Sys\_fork errors

EAGAIN

ENOMEM

Too many processes already exist. Sufficient virtual memory for the new process was not available.

## Sys\_execv

- Replace the currently executing program image with a new process image.
- Process id is unchanged.
- int sys\_execv(char \*program, char \*\*args)
  - program: path name of the program to run.
  - Args: tf->tf\_a0 and tf->tf\_a1
- Most of the implementation is already in the runprogram.c
  - Only a few more things.
  - Check the last argument in \*\*args is NULL.
  - Make sure it is less than MAX\_ARGS\_NUM
  - *copyin* the arguments from user space to kernel space.
  - Create a new address space.
    - as\_create()
  - Allocate a stack on it.
    - as\_define\_stack()
  - Copyout the arguments back onto the new stack

### Sys\_execv errors

ENODEV	The device prefix of program did not exist.
ENOTDIR	A non-final component of <i>program</i> was not a directory.
ENOENT	<i>program</i> did not exist.
EISDIR	program is a directory.
ENOEXEC	<i>program</i> is not in a recognizable executable file format, was for the wrong platform, or contained invalid fields.
ENOMEM	Insufficient virtual memory is available.
E2BIG	The total size of the argument strings is too large.
EIO	A hard I/O error occurred.
EFAULT	One of the args is an invalid pointer.

## Sys\_waitpid

- Wait for the process with pid to exit.
- Return its exit code via the integer pointer status.
- pid\_t sys\_waitpid(pid\_t pid, int \*status, int options)
- You need a mechanism for processes to show *interest* into each other.
  - •You can add restrictions on which processes can show interest.
  - •Make sure to prevent deadlocks by either setting restrictions to prevent it or to implement a mechanism to detect it.
- Return the pid with status assigned to exit status on success.
- If error, return -1 and set the ret pointer to the error code.

## Sys\_waitpid errors

EINVAL

EFAULT

The *options* argument requested invalid or unsupported options.

The *status* argument was an invalid pointer.

## Sys\_exit

- Causes the current process to terminate.
- The process id of the exiting process cannot be reused if there are other processes *interested* in it.
  - Do not put the exited pid back to available pid pool blindly.
- void sys\_\_\_exit(int code)
  - Code is the exitcode that will be given to other processes who are *interested* in it

## Scheduler

- Currently os161 has single queue round-robin scheduler.
- You can modify hardclock.c to have another counter that counts in HZ/2.
- Mostly scheduler.c will be edited.
  - Add a new queue.
  - Add each process a priority and modify make\_runnable to match the thread and queue level according to its priority.
  - Modify the scheduler function such that the chances of picking higher level queue will increase.

## Testing

- os161/man/testbin has the details about given tests
  - Contains html files
  - Read them carefully and understand what needs to be implemented to pass the tests
  - Be careful: some of them requires VM management to work
- Forktest is very useful
- Also test bin/cp example in the assignment description
- Shell implementation is given but not necessary
  - You can call the tests by p /testbin/forktest
- A basic sys\_write is also provided. It will be necessary for printf statements from inside a user-program

## Testing

- Build you own tests
- Repeat some of the tests with your new scheduler enabled
  Report the response times with different quantum sizes
- Make sure to include all the test outputs in your submission

## Thank you