

MapReduce

CS 475: Concurrent & Distributed Systems (Fall 2021) Lecture 4

Yue Cheng

Some material taken/derived from:

• Princeton COS-418 materials created by Michael Freedman and Wyatt Lloyd.

• MIT 6.824 by Robert Morris, Frans Kaashoek, and Nickolai Zeldovich.

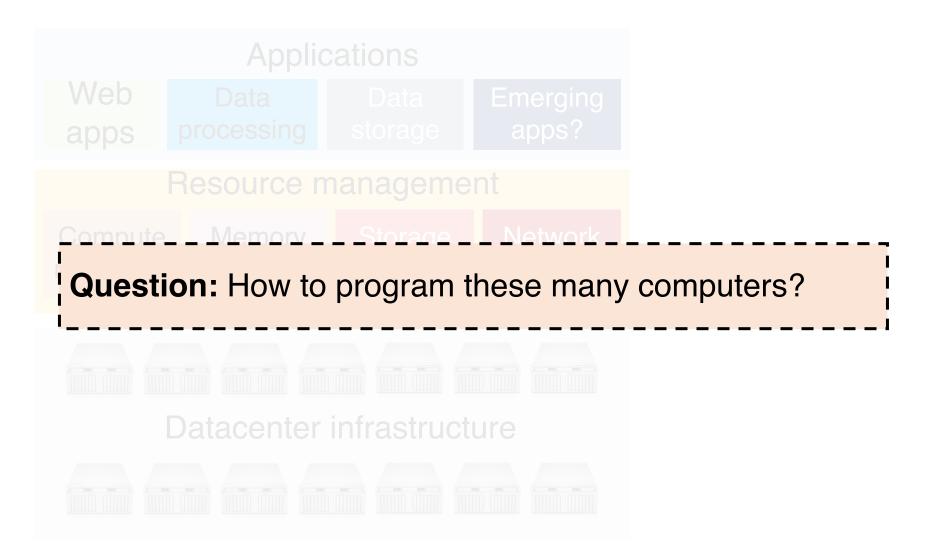
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Applications				
Web	Data	Data	Emerging	
apps	processing	storage	apps?	
Resource management				
Compute resources	Memory resources	Storage resources	Network resources	

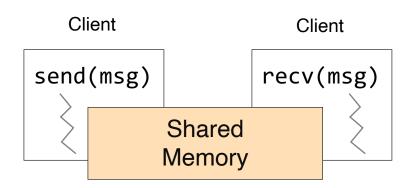


Datacenter infrastructure



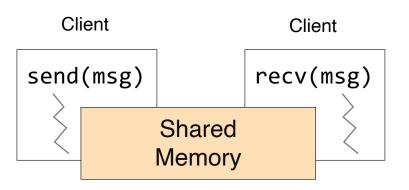


Review: Shared memory

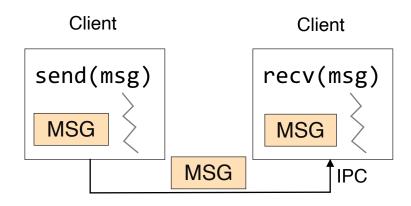


- Shared memory: multiple processes to share data via memory
- Applications must locate and and map shared memory regions to exchange data

Review: Shared memory vs. Message passing



- Shared memory: multiple processes to share data via memory
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- Message passing: exchange data explicitly via IPC
- Application developers define protocol and exchanging format, number of participants, and each exchange

Review: Shared memory vs. Message passing

- Easy to program; just like a single multithreaded machines
- Hard to write high perf. apps:
 - Cannot control which data is local or remote (remote mem. access much slower)
- Hard to mask failures

- Message passing: can write very high perf. apps
- Hard to write apps:
 - Need to manually decompose the app, and move data
- Need to manually handle failures

Shared memory: Pthread

- A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization
- API specifies behavior of the thread library, implementation is up to development of the library
- Common in UNIX (e.g., Linux) OSes

Shared memory: Pthread

```
void *myThreadFun(void *vargp) {
    sleep(1);
    printf("Hello world\n");
    return NULL;
}
int main() {
    pthread_t thread_id_1, thread_id_2;
    pthread_create(&thread_id_1, NULL, myThreadFun, NULL);
    pthread_create(&thread_id_2, NULL, myThreadFun, NULL);
    pthread_join(thread_id_1, NULL);
    pthread_join(thread_id_2, NULL);
    exit(0);
}
```

Message passing: MPI

- MPI Message Passing Interface
 - Library standard defined by a committee of vendors, implementers, and parallel programmers
 - Used to create parallel programs based on message passing
- Portable: one standard, many implementations
 - Available on almost all parallel machines in C and Fortran
 - De facto standard platform for the HPC community

Message passing: MPI

```
int main(int argc, char **argv) {
      MPI Init(NULL, NULL);
      // Get the number of processes
      int world_size;
      MPI_Comm_size(MPI_COMM_WORLD, &world_size);
      // Get the rank of the process
      int world rank;
      MPI Comm rank(MPI COMM WORLD, *world rank);
      // Print off a hello world message
      printf("Hello world from rank %d out of %d processors\n",
            world rank, workld size);
      // Finalize the MPI environment
      MPI Finalize();
}
```

Message passing: MPI

mpirun -n 4 -f host_file ./mpi_hello_world

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}

MapReduce

The big picture (motivation)

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- Datasets are too big to process using a single computer
- Good parallel processing engines are rare (back then in the late 90s)
- Want a parallel processing framework that:
 - is general (works for many problems)
 - is easy to use (no locks, no need to explicitly handle communication, no race conditions)
 - can automatically parallelize tasks
 - can automatically handle machine failures

Context (Google circa 2000)

- Starting to deal with massive datasets
- But also addicted to cheap, unreliable hardware
 - Young company, expensive hardware not practical
- Only a few expert programmers can write distributed programs to process them
 - Scale so large jobs can complete before failures



Context (Google circa 2000)

- Starting to deal with massive datasets
- But also addicted to cheap, unreliable hardware
 - Young company, expensive hardware not practical
- Only a few expert programmers can write distributed programs to process them
 - Scale so large jobs can complete before failures
- Key question: how can every Google engineer be imbued with the ability to write parallel, scalable, distributed, fault-tolerant code?
- Solution: abstract out the redundant parts
- Restriction: relies on job semantics, so restricts which problems it works for

Application: Word Count

SELECT count(word), word FROM data GROUP BY word

Deal with multiple files?

1. Compute word counts from individual files

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- 2. Then merge intermediate output

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- 1. Compute word counts from individual files
- 2. Then merge intermediate output
- 3. Compute word count on merged outputs

What if the data is too big to fit in one computer?

- 1. In parallel, send to worker:
 - Compute word counts from individual files
 - Collect results, wait until all finished

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MapReduce: Programming interface

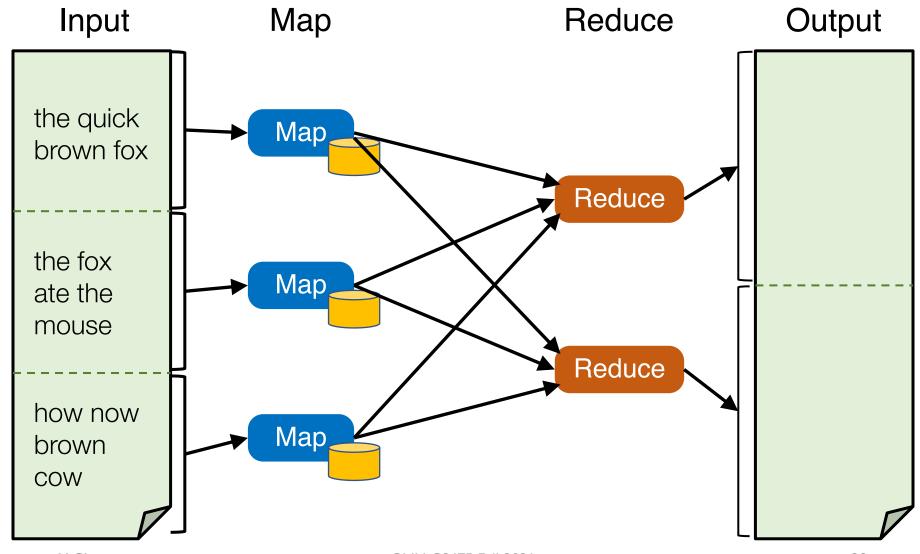
- map(k1, v1) \rightarrow list(k2, v2)
 - Apply function to (k1, v1) pair and produce set of intermediate pairs (k2, v2)

- reduce(k2, list(v2)) \rightarrow list(k3, v3)
 - Apply aggregation (reduce) function to values
 - Output results

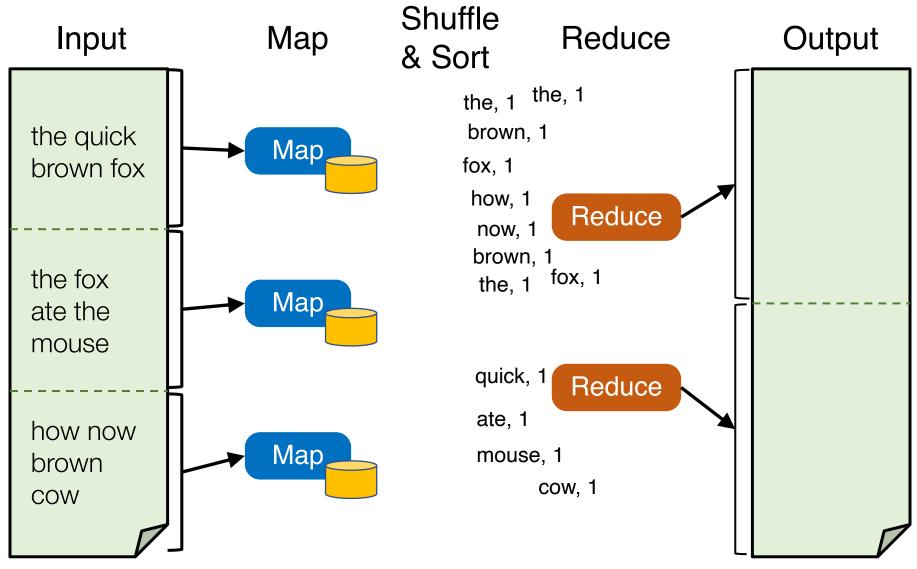
MapReduce: Word Count

```
map(key, value):
   for each word w in value:
       EmitIntermediate(w, "1");
reduce(key, values):
   int result = 0;
   for each v in values:
       results += ParseInt(v);
   Emit(AsString(result));
```

Word Count execution



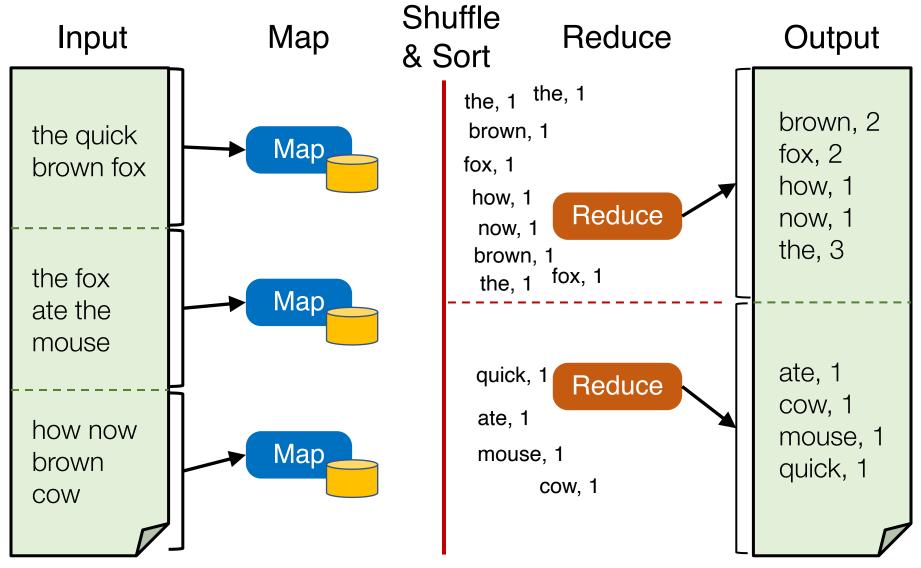
Word Count execution



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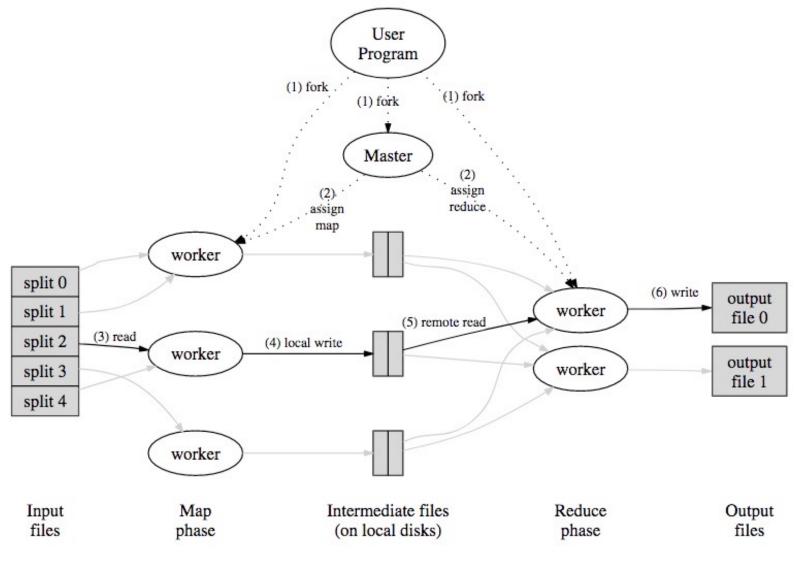
Word Count execution



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MapReduce data flows



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MapReduce processes

Map Shuffle & Sort Reduce

Reduce

Reduce

- Map workers write intermediate output to local disk, separated by partitioning. Once completed, tell master node
- Reduce worker told of location of map task outputs, pulls their partition's data from each mapper, execute function across data
- Note:
 - "All-to-all" shuffle b/w mappers and reducers
 - Written to disk ("materialized") b/w each state

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Map

Map

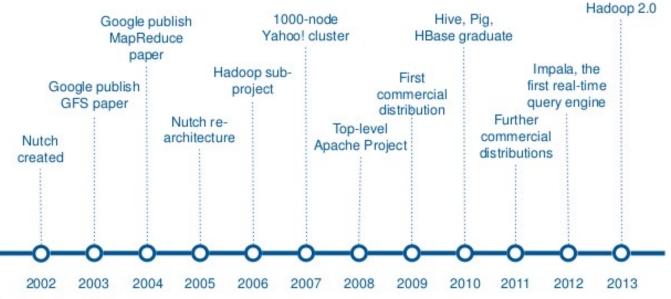
Map

Apache Hadoop



- An open-source implementation of Google's MapReduce framework
 - Hadoop MapReduce atop Hadoop Distributed File System (HDFS)

A Brief History of Hadoop



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