MapReduce

CS 4740: Cloud Computing Fall 2024 Lecture 4

Yue Cheng



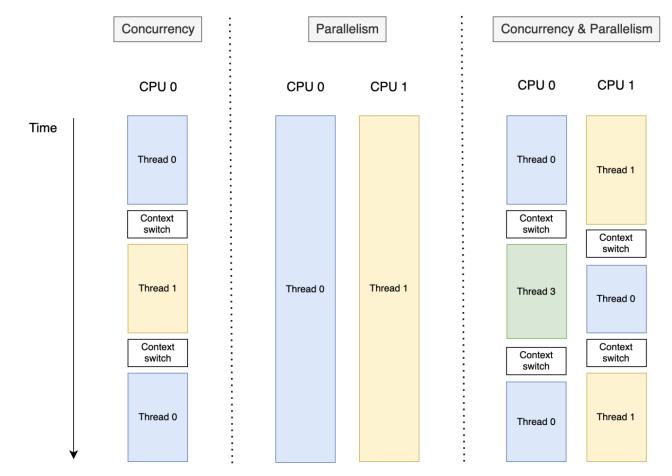
Some material taken/derived from:

- Princeton COS-418 materials created by Michael Freedman and Wyatt Lloyd.
- MIT 6.824 by Robert Morris, Frans Kaashoek, and Nickolai Zeldovich.

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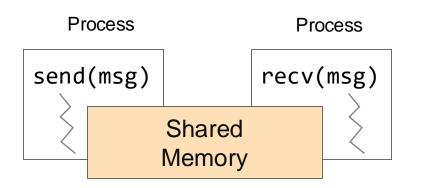
Recap: Parallelism vs. concurrency

"Concurrency is about dealing with lots of things at once. Parallelism is about doing lots of things at once."



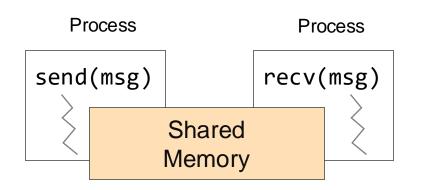
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Recap: Shared memory

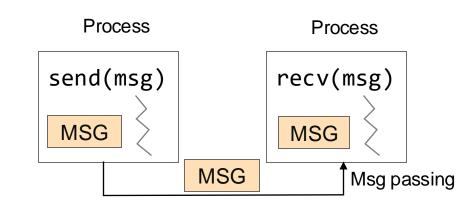


- Shared memory: multiple processes to share data via memory
- Applications must locate and and map shared memory regions to exchange data

Recap: Shared memory vs. Message passing



- Shared memory: multiple processes to share data via memory
- Applications must locate and and map shared memory regions to exchange data



- Message passing: exchange data explicitly via message passing
- Application developers define protocol and exchanging format, number of participants, and each exchange

Recap: Shared memory vs. Message passing

- Easy to program; just like a single multithreaded machines
- Hard to write highperformance apps:
 - Cannot control which data is local or remote (remote mem. access much slower)
- Hard to mask failures

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- Easy to program; just like a single multithreaded machines
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- Message passing: can write very highperformance apps
- Hard to write apps:
 - Need to manually decompose the app, and move data
- Need to manually handle failures

Shared memory: Pthread

- A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization
- API specifies behavior of the thread library, implementation is up to development of the library
- Common in UNIX (e.g., Linux) OSes

Shared memory: Pthread

```
void *myThreadFun(void *vargp) {
    sleep(1);
    printf("Hello world\n");
    return NULL;
}
int main() {
    pthread_t thread_id_1, thread_id_2;
    pthread_create(&thread_id_1, NULL, myThreadFun, NULL);
    pthread_create(&thread_id_2, NULL, myThreadFun, NULL);
    pthread_join(thread_id_1, NULL);
    pthread_join(thread_id_2, NULL);
    exit(0);
}
```

Message passing: MPI

- MPI Message Passing Interface
 - Library standard defined by a committee of vendors, implementers, and parallel programmers
 - Used to create parallel programs based on message passing
- Portable: one standard, many implementations
 - Available on almost all parallel machines in C and Fortran
 - De facto standard platform for the HPC community

Message passing: MPI

```
int main(int argc, char **argv) {
      MPI Init(NULL, NULL);
      // Get the number of processes
      int world size;
      MPI Comm size(MPI COMM WORLD, &world size);
      // Get the rank of the process
      int world_rank;
      MPI Comm rank(MPI COMM WORLD, *world rank);
      // Print off a hello world message
      printf("Hello world from rank %d out of %d processors\n",
            world rank, world size);
      // Finalize the MPI environment
      MPI Finalize();
}
```

Message passing: MPI

mpirun -n 4 -f host_file ./mpi_hello_world

```
int main(int argc, char **argv) {
    MPI_Init(NULL, NULL);
```

```
// Get the number of processes
int world_size;
MPI_Comm_size(MPI_COMM_WORLD, &world_size);
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}

MapReduce

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- Good parallel processing engines are rare (back then in the late 90s)
- Want a parallel processing framework that:
 - is **general** (works for many problems)
 - is **easy to use** (no locks, no need to explicitly handle communication, no race conditions)
 - can automatically parallelize tasks
 - can automatically handle machine failures

Context (Google circa 2000)

- Starting to deal with massive datasets
- But also addicted to cheap, commodity hardware
 - Young company, expensive hardware not practical
- Only a few expert programmers can write distributed programs to process them
 - Scale so large jobs can complete despite failures



Context (Google circa 2000)

- Starting to deal with massive datasets
- But also addicted to cheap, commodity hardware
 - Young company, expensive hardware not practical
- Only a few expert programmers can write distributed programs to process them
 - Scale so large jobs can complete despite failures
- Key question: how can every Google engineer be imbued with the ability to write parallel, scalable, distributed, fault-tolerant code?
- Solution: abstract out the redundant parts
- Restriction: relies on job semantics, so restricts which problems it works for

Application: Word Count

SELECT count(word), word FROM data GROUP BY word

Deal with multiple files?

1. Compute word counts from individual files

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 - Compute word counts from individual files
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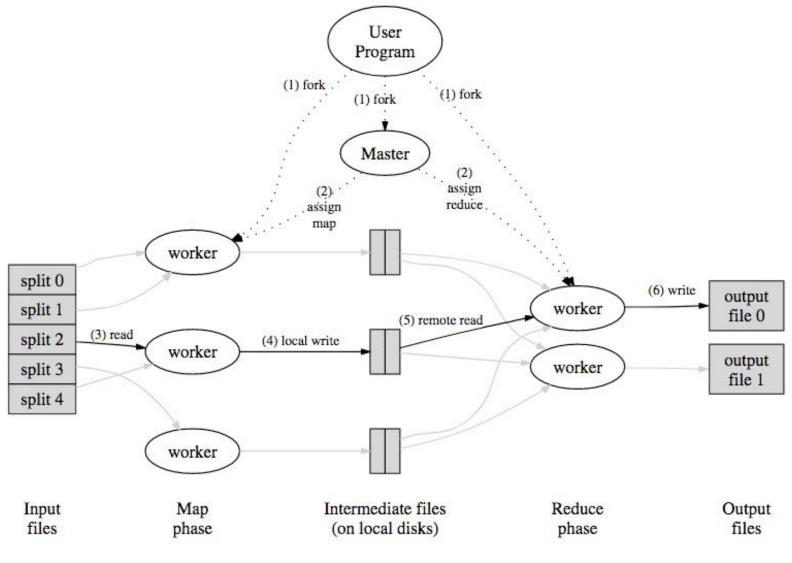
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MapReduce: Programming interface

- map(k1, v1) \rightarrow list(k2, v2)
 - Apply function to (k1, v1) pair and produce set of intermediate pairs (k2, v2)

- reduce(k2, list(v2)) \rightarrow list(k3, v3)
 - Apply aggregation (reduce) function to values
 - Output results

MapReduce data flows



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MapReduce visualization

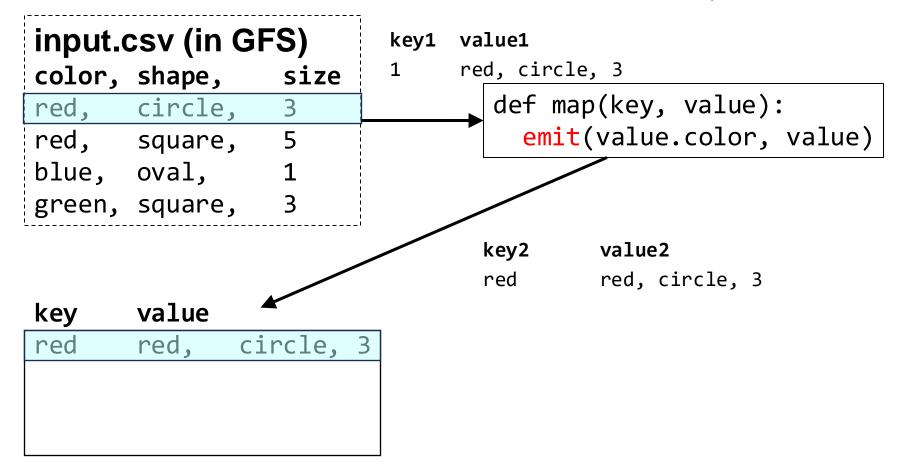
input.csv (in GFS)			
color,	shape,	size	
red,	circle,	3	
red,	square,	5	
blue,	oval,	1	
green,	square,	3	

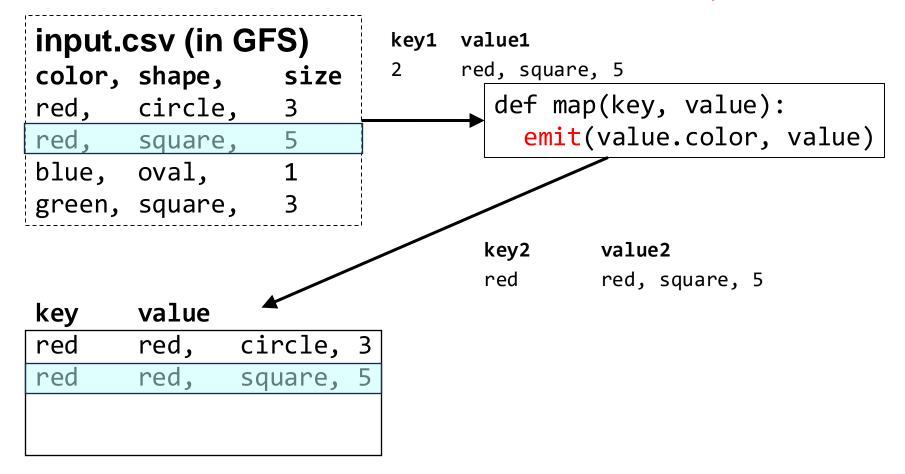
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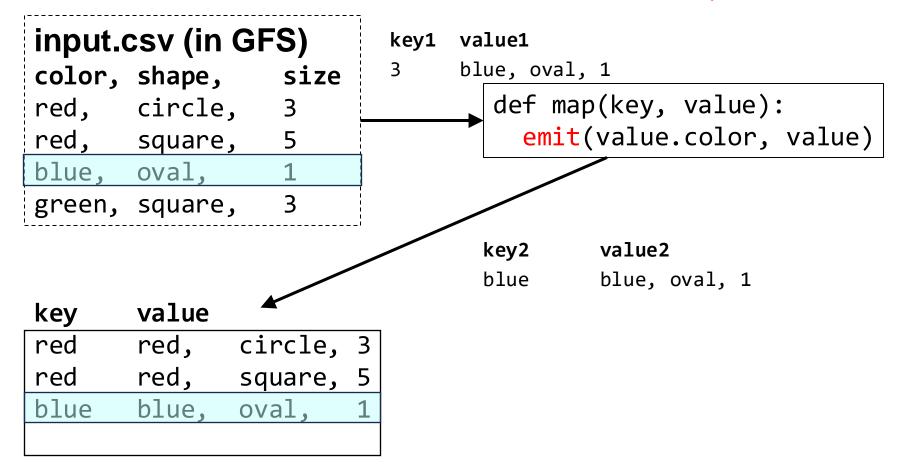
How to count the number of occurrences for each unique color?

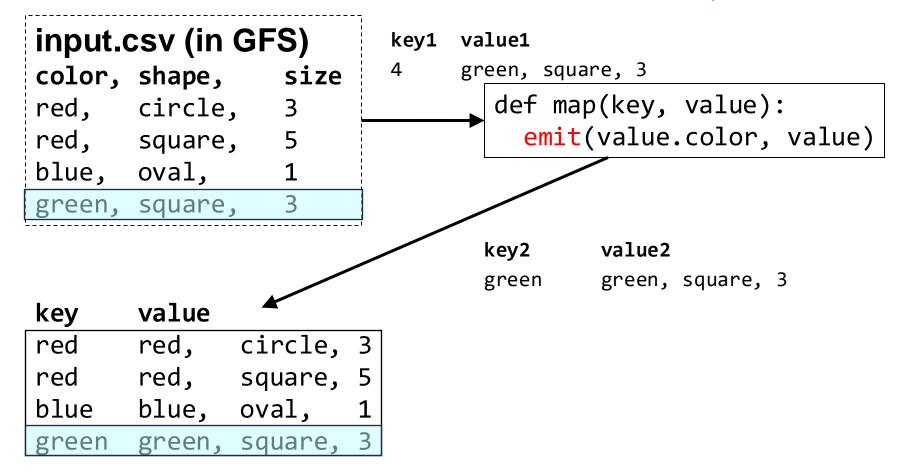
def map(key, value):
 emit(value.color, value)

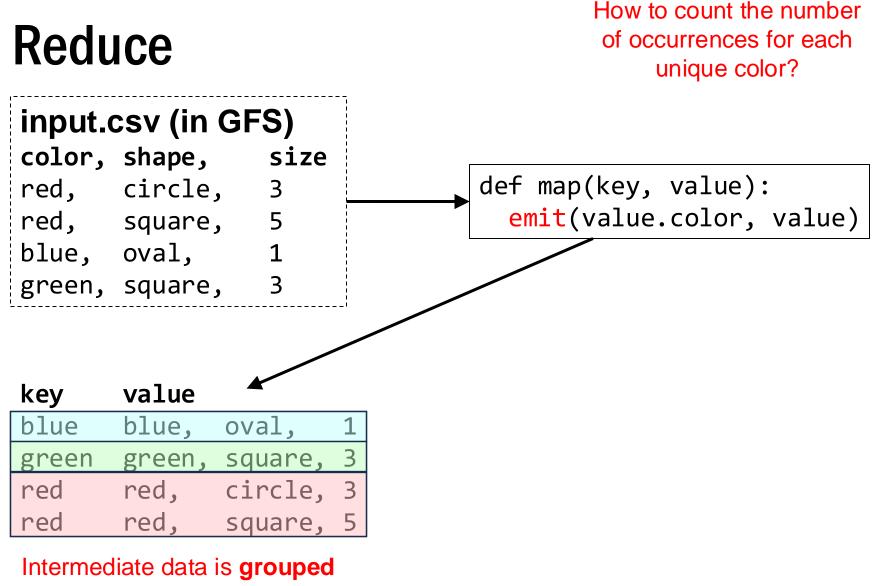
Map will be called 4 times (once for each line of the input file).



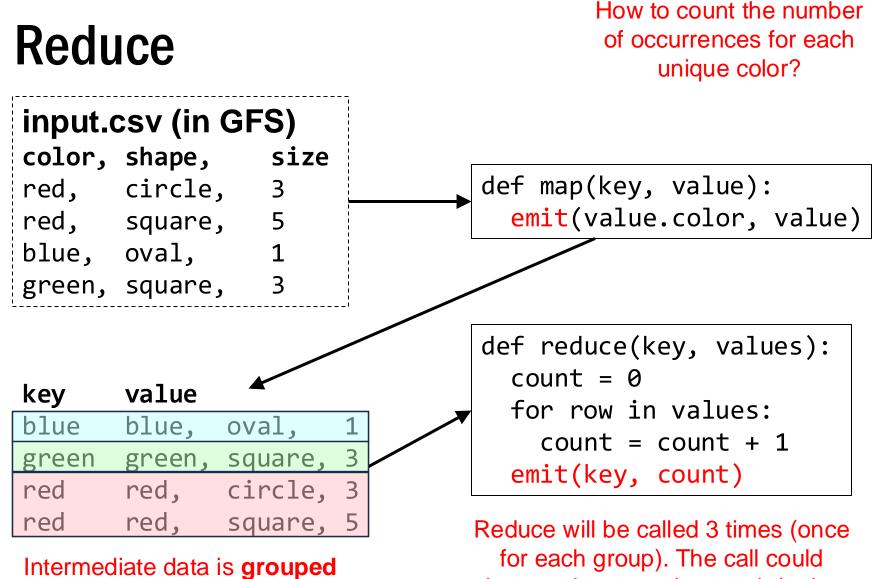




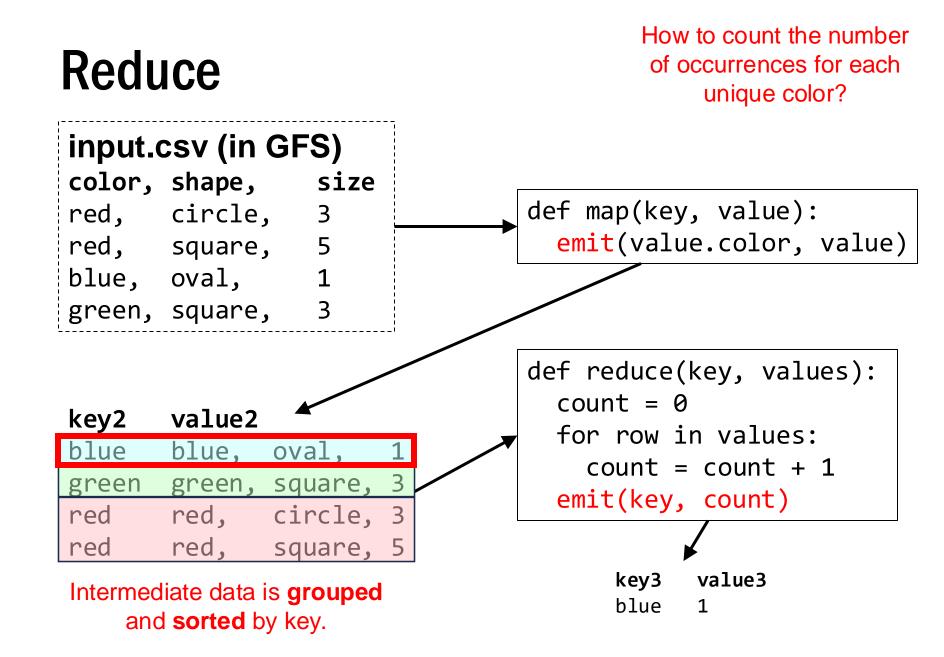


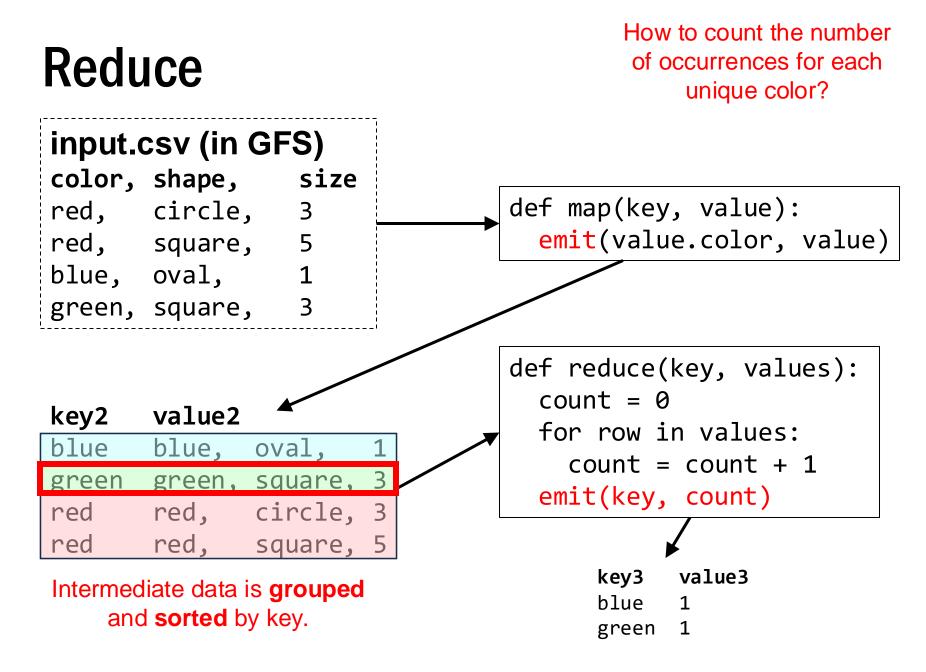


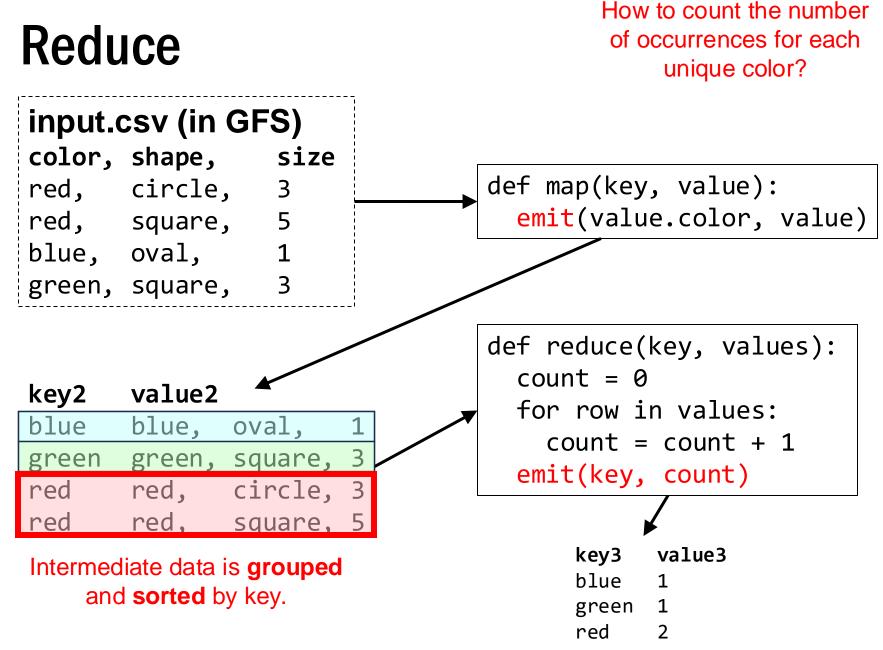
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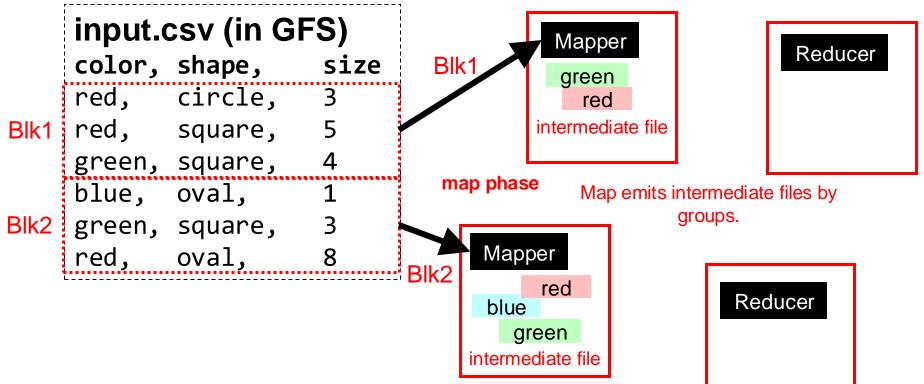






Multiple reducers (for big intermediate data)

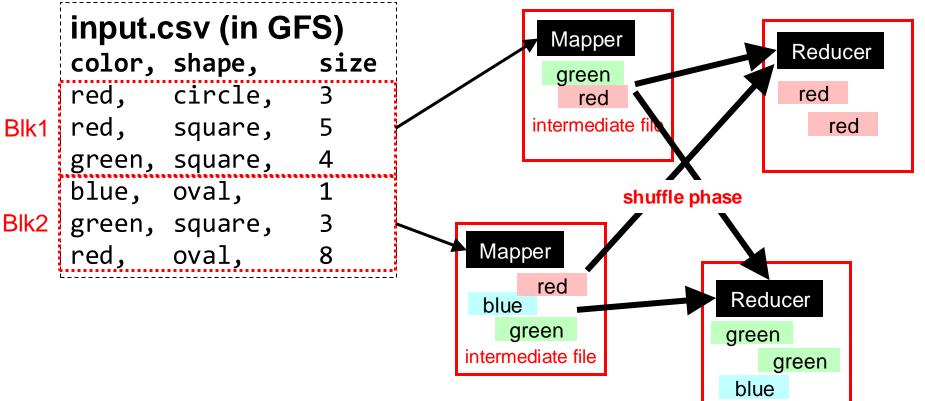
Cluster of machines



Each reduce task produces one output file. A reduce task might take multiple keys. All intermediate rows with the same key go to the same reducer.

Multiple reducers (for big intermediate data)

Cluster of machines

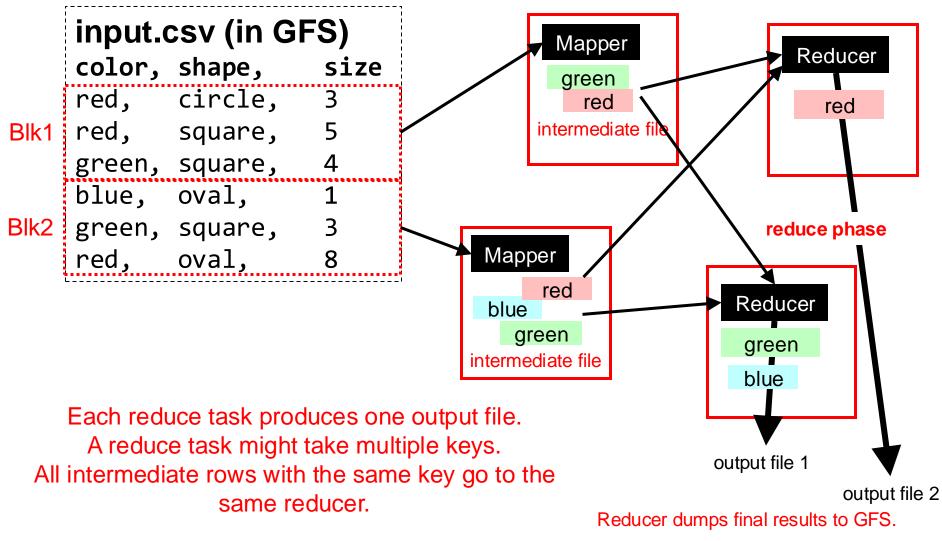


Each reduce task produces one output file.

A reduce task might take multiple keys. All intermediate rows with the same key go to the same reducer. Reducer collects all intermediate files of its assigned keys (groups).

Multiple reducers (for big intermediate data)

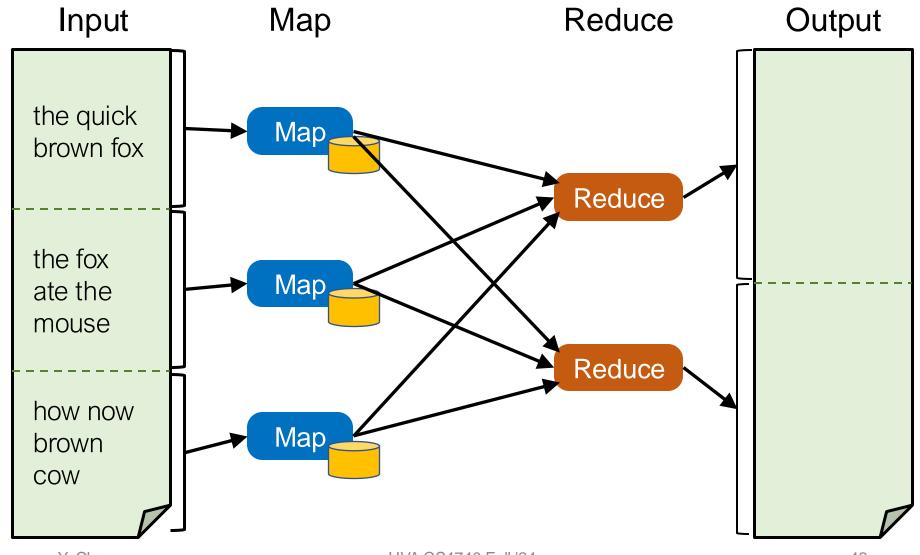
Cluster of machines



MapReduce: Word Count

```
map(key, value):
   for each word w in value:
       EmitIntermediate(w, "1");
reduce(key, values):
   int result = 0;
   for each v in values:
       results += ParseInt(v);
   Emit(AsString(result));
```

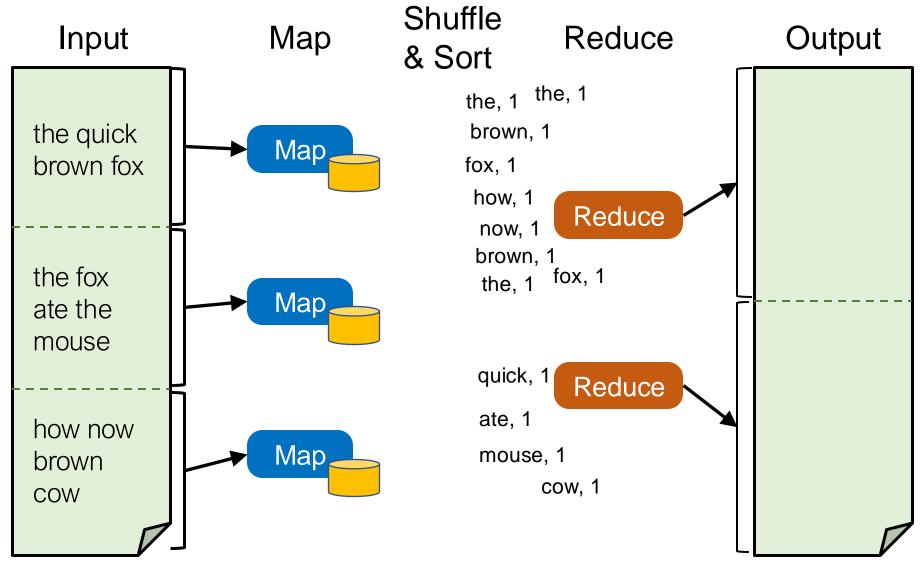
Word Count execution



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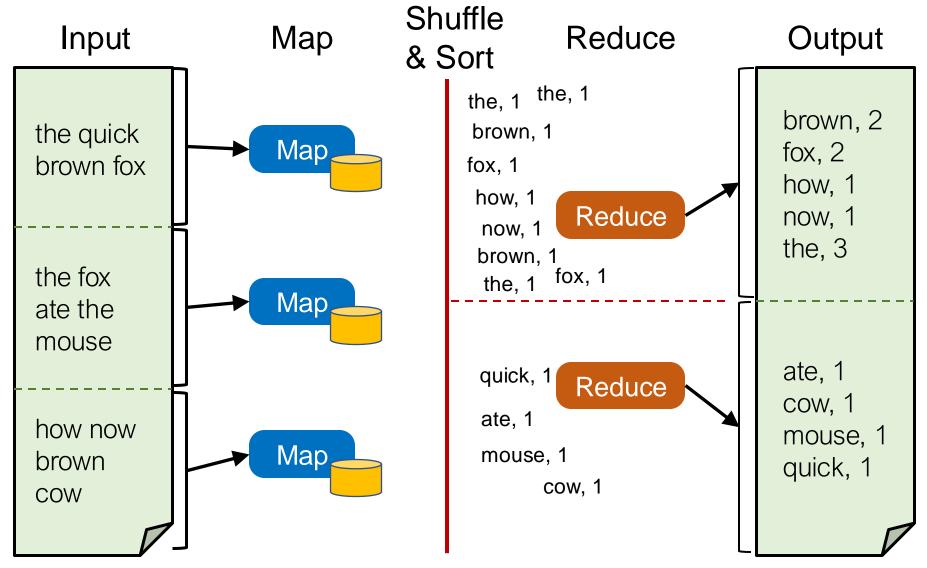
Word Count execution



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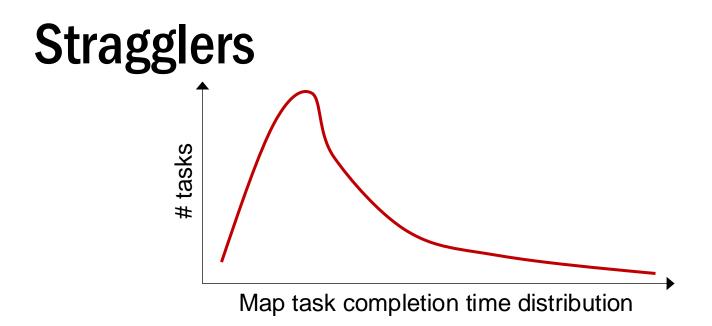
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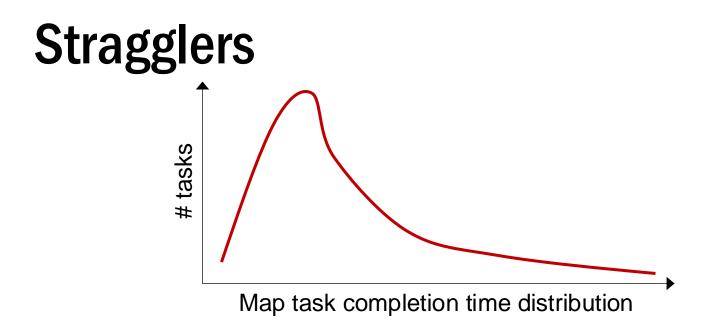
Word Count execution



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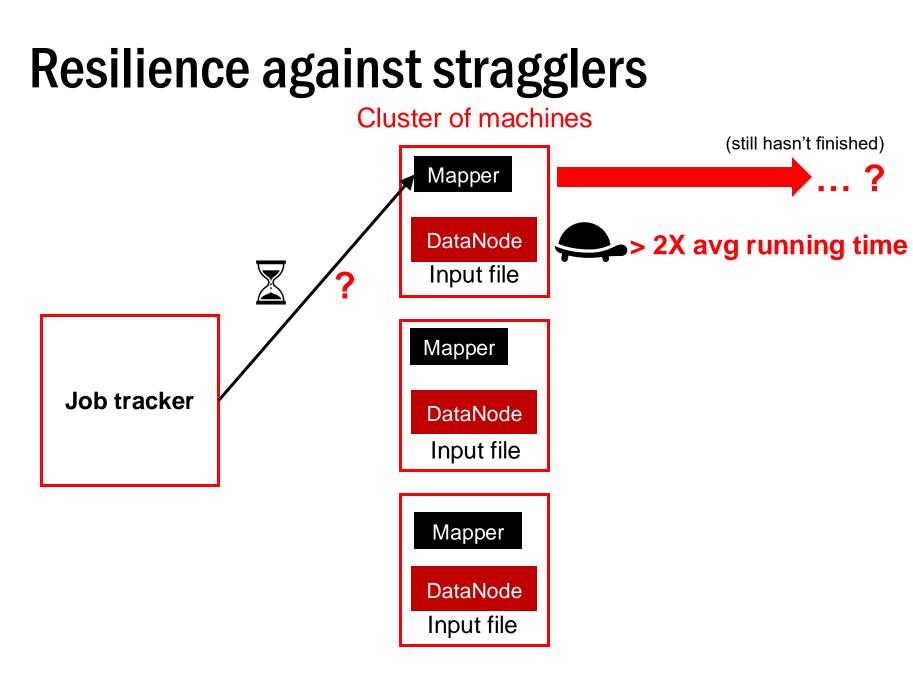
• Tail execution time means some executors (always) finish late (tail latency)

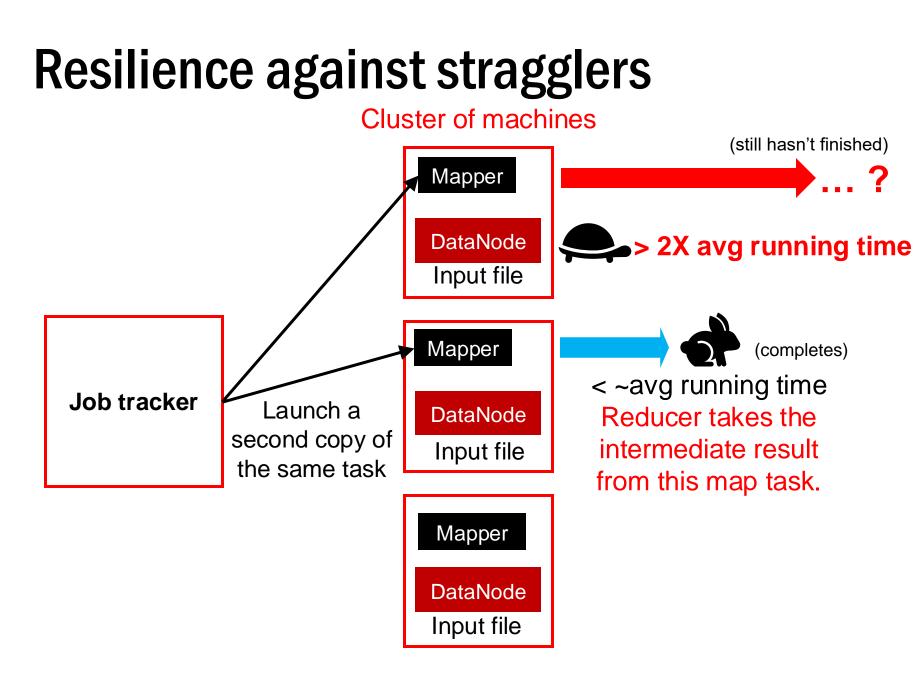
Q: How can MapReduce work around this?

Hint: its approach to fault-tolerance provides the right tool

Resilience against stragglers?

- If a task is going slowly (i.e., straggler):
 - Launch second copy of task (backup task) on another node
 - Take the output of whichever finishes first





Would backup tasks cause correctness issue in MapReduce jobs?

Discussion: MapReduce eval (paper)

