MapReduce

CS 4740: Cloud Computing Fall 2024 Lecture 4

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Some material taken/derived from:

- Princeton COS-418 materials created by Michael Freedman and Wyatt Lloyd.
- MIT 6.824 by Robert Morris, Frans Kaashoek, and Nickolai Zeldovich.

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Recap: Parallelism vs. concurrency

"Concurrency is about dealing with lots of things at once. Parallelism is about doing lots of things at once."

Recap: Shared memory

- Shared memory: multiple processes to share data via memory
- Applications must locate and and map shared memory regions to exchange data

Recap: Shared memory vs. Message passing

- Shared memory: multiple processes to share data via memory
- Applications must locate and and map shared memory regions to exchange data

- Message passing: exchange data explicitly via message passing
- Application developers define protocol and exchanging format, number of participants, and each exchange

Recap: Shared memory vs. Message passing

- Easy to program; just like a single multithreaded machines
- Hard to write highperformance apps:
	- Cannot control which data is local or remote (remote mem. access much slower)
- Hard to mask failures

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- Easy to program; just like a single multithreaded machines
- Hard to write highperformance apps:
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- Hard to mask failures
- Message passing: can write very highperformance apps
- Hard to write apps:
	- Need to manually decompose the app, and move data
- Need to manually handle failures

Shared memory: Pthread

- A POSIX standard (IEEE 1003.1c) API for thread creation and synchronization
- API specifies behavior of the thread library, implementation is up to development of the library
- Common in UNIX (e.g., Linux) OSes

Shared memory: Pthread

```
void *myThreadFun(void *vargp) {
      sleep(1);
      printf("Hello world\n");
      return NULL;
}
int main() {
      pthread_t thread_id_1, thread_id_2;
      pthread_create(&thread_id_1, NULL, myThreadFun, NULL);
      pthread create(&thread id 2, NULL, myThreadFun, NULL);
      pthread join(thread id 1, NULL);
      pthread join(thread id 2, NULL);
      exit(0);}
```
Message passing: MPI

- MPI Message Passing Interface
	- Library standard defined by a committee of vendors, implementers, and parallel programmers
	- Used to create parallel programs based on message passing
- Portable: one standard, many implementations
	- Available on almost all parallel machines in C and Fortran
	- De facto standard platform for the HPC community

Message passing: MPI

```
int main(int argc, char **argv) {
      MPI Init(NULL, NULL);
      // Get the number of processes
      int world_size;
      MPI_Comm_size(MPI_COMM_WORLD, &world_size);
      // Get the rank of the process
      int world_rank;
      MPI Comm rank(MPI_COMM_WORLD, *world_rank);
      // Print off a hello world message
      printf("Hello world from rank %d out of %d processors\n",
            world rank, world size);
      // Finalize the MPI environment
      MPI Finalize();
}
```
Message passing: MPI

mpirun –n 4 –f host_file ./mpi_hello_world

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MapReduce

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- Datasets are too big to process using a single computer
- Good parallel processing engines are rare (back then in the late 90s)
- Want a parallel processing framework that:
	- is general (works for many problems)
	- is easy to use (no locks, no need to explicitly handle communication, no race conditions)
	- can automatically parallelize tasks
	- can automatically handle machine failures

Context (Google circa 2000)

- Starting to deal with massive datasets
- But also addicted to cheap, commodity hardware
	- Young company, expensive hardware not practical
- Only a few expert programmers can write distributed programs to process them
	- Scale so large jobs can complete despite failures

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- Only a few expert programmers can write distributed programs to process them
	- Scale so large jobs can complete despite failures
- Key question: how can every Google engineer be imbued with the ability to write parallel, scalable, distributed, fault-tolerant code?
- Solution: abstract out the redundant parts
- Restriction: relies on job semantics, so restricts which problems it works for

Application: Word Count

```
cat data.txt
     | tr -s '[[:punct:][:space:]]' '\n'
     | sort | uniq -c
```
SELECT count(word), word FROM data GROUP BY word

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MapReduce: Programming interface

- map(k1, v1) \rightarrow list(k2, v2)
	- Apply function to (k1, v1) pair and produce set of intermediate pairs (k2, v2)

- reduce(k2, list(v2)) \rightarrow list(k3, v3)
	- Apply aggregation (reduce) function to values
	- Output results

MapReduce data flows

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MapReduce visualization

input.csv (in GFS) color, shape, size red, circle, 3 red, square, 5 blue, oval, 1 green, square, 3

How to count the number of occurrences for each unique color?

def map(key, value): emit(value.color, value)

Map will be called 4 times (once for each line of the input file).

and **sorted** by key.

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Multiple reducers (for big intermediate data)

Cluster of machines

Each reduce task produces one output file. A reduce task might take multiple keys. All intermediate rows with the same key go to the same reducer.

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Reducer collects all intermediate files of its assigned keys (groups).

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Cluster of machines

MapReduce: Word Count

```
map(key, value):
   for each word w in value:
       EmitIntermediate(w, "1");
reduce(key, values):
   int result = 0;
   for each v in values:
       results += ParseInt(v);
   Emit(AsString(result));
```
Word Count execution

Word Count execution

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Word Count execution

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- **Tail execution time** means some executors (always) finish late (tail latency)
- Q: How can MapReduce work around this?
	- Hint: its approach to **fault-tolerance** provides the right tool

Resilience against stragglers?

- If a task is going slowly (i.e., **straggler**):
	- Launch second copy of task (**backup task**) on another node
	- Take the output of whichever finishes first

Would backup tasks cause correctness issue in MapReduce jobs?

Discussion: MapReduce eval (paper)

